
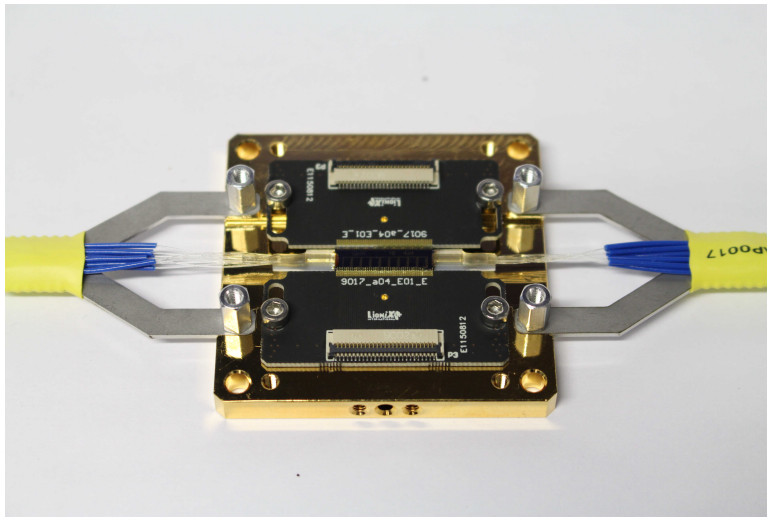
 **Boson sampling simulation  
enhanced by FPGA based  
data-flow engines**

 **WIGNER**

Peter Rakyta

Department of Physics of  
Complex Systems

## optical interferometers on chip



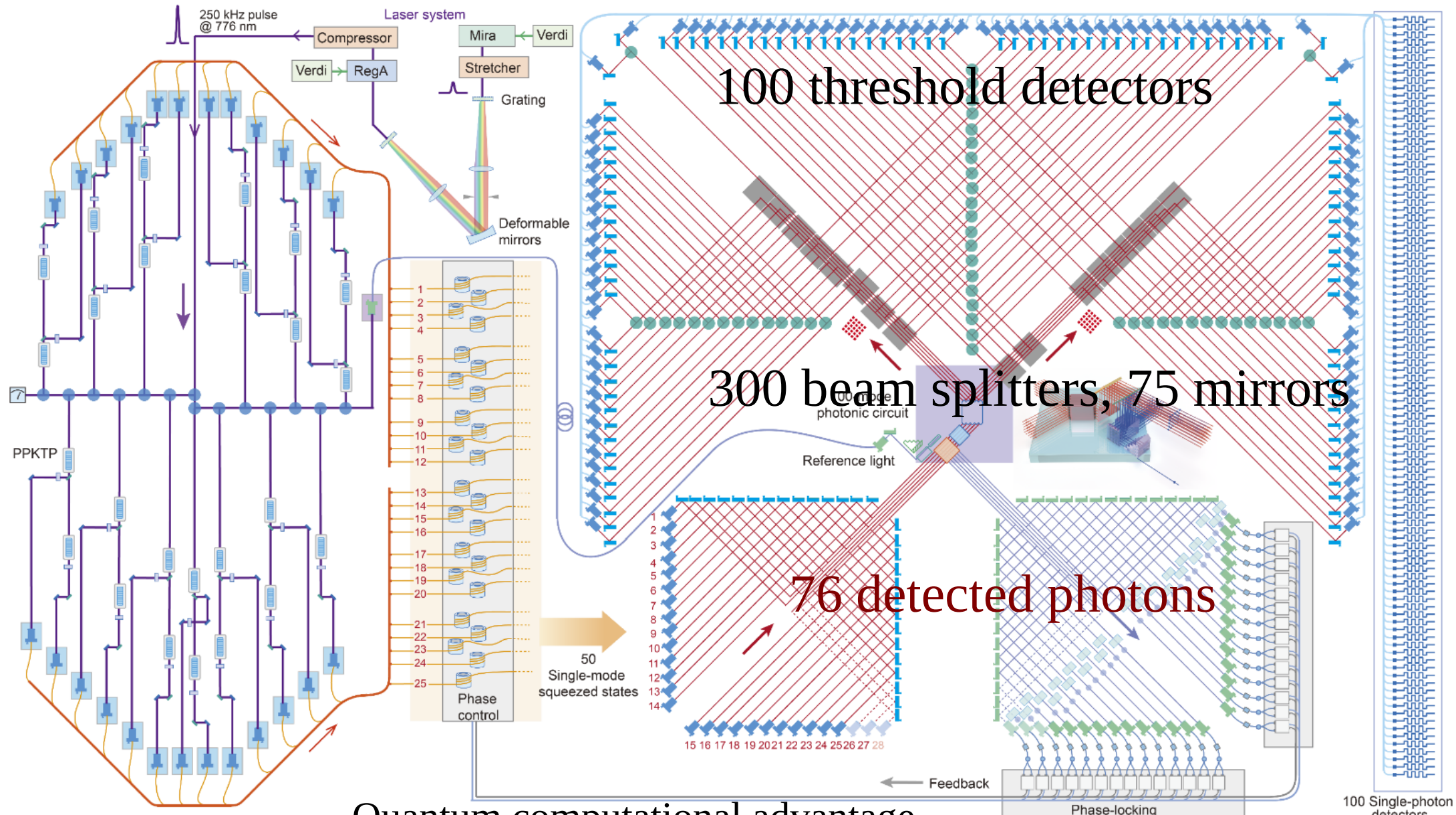
- N optical modes occupied by 0,1,2, ... N photons (instead of qubits)
- single photon sources or Gaussian states

generated from coherent laser beams

on optical table



# Quantum advantage showed on Boson Sampling from a Gaussian state



Quantum computational advantage using photons. *Science*, 370(6523):1460–1463, 2020.

University of Science and Technology of China  
 Chinese Academy of Sciences, Tsinghua University, China



## Spin-off University of Twente



THE FASTEST WAY TO A QUANTUM FUTURE



## Netherlands

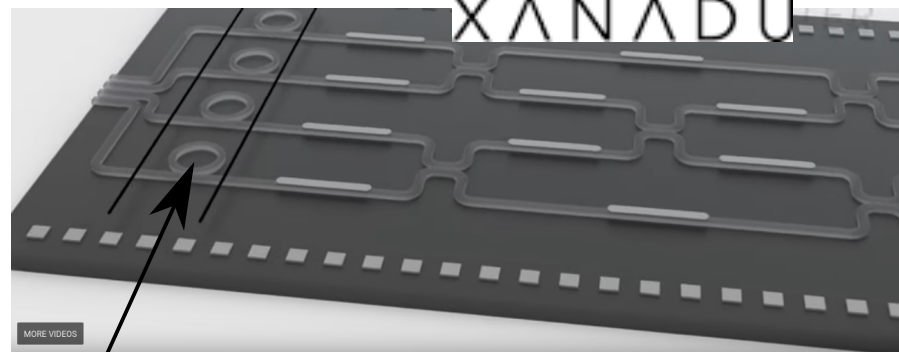
connected to single photon sources and PNR detectors

12 photonic modes

## Canada



XANADU



creates squeezed state from classical laser input - this is the qubit

24 photonic modes

### Strawberry Fields:

quantum computer simulator

### PennyLane:

first library for quantum machine learning



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# The Piquasso project

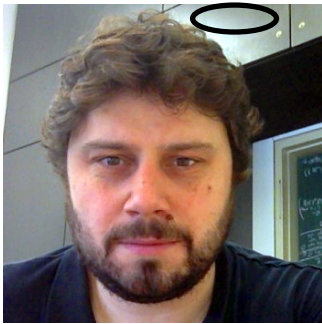


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PIQUASSO

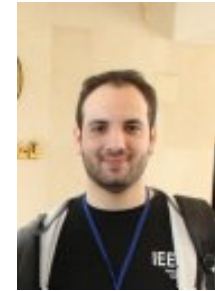
Why don't we put together a bosonic quantum simulator?



Zoltán Zimborás  
Wigner



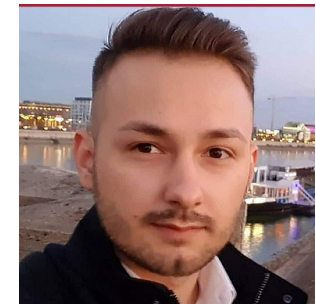
Poór Boldizsár



Kareem El Safty



Kozsik tamás



Jóczik Szabolcs



Michał Oszmaniec



Tomasz Rybotycki



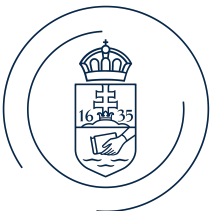
Kaposi Ágoston



Kolarovszki Zoltán



Supported by



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# Boson Sampling



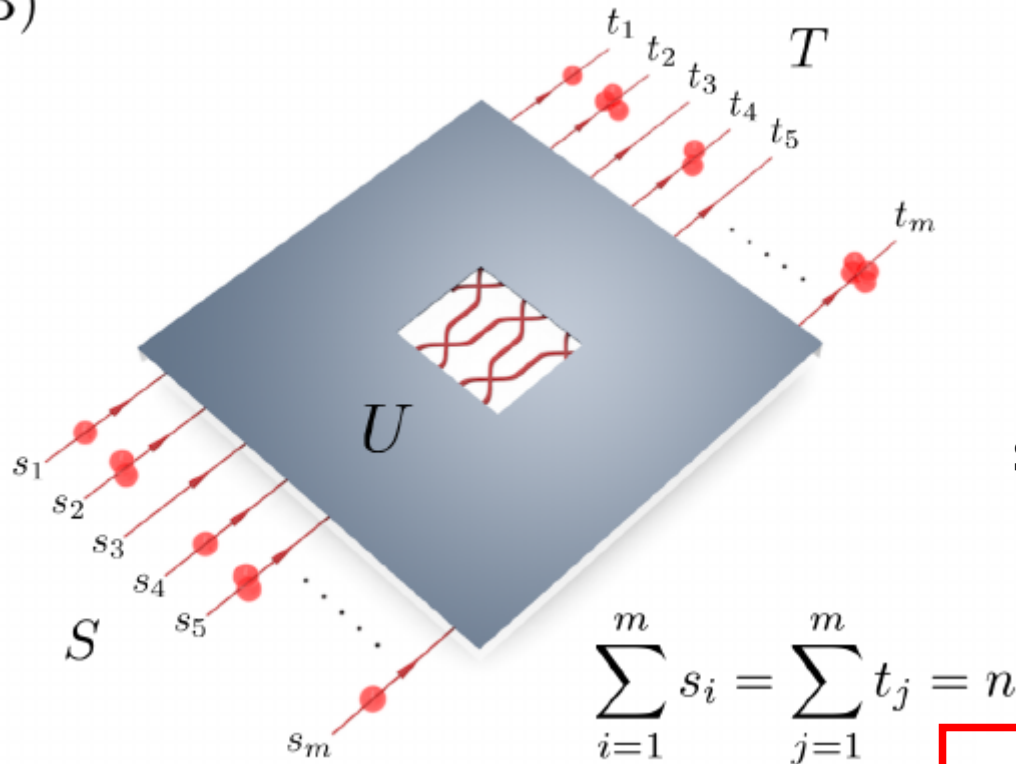
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## The Computational Complexity of Linear Optics

Scott Aaronson\*

Alex Arkhipov†

(B)



$$P[S \rightarrow T] = \frac{|Per(U_{S,T})|^2}{s_1! \dots s_m! t_1! \dots t_m!}$$

$$Per(A) = \sum_{\sigma \in S_n} \prod_{i=1}^n a_{i\sigma(i)}$$

Set of permutations of (1,2,3 ... n)

$$\mathcal{O}(n^2 \cdot 2^n)$$

$$\mathcal{O}((n-1) \cdot n!)$$

Ryser's formula:

$$Per(A) = (-1)^n \sum_{S \subseteq \{1,2,\dots,n\}} (-1)^{|S|} \prod_{i=1}^n \sum_{j \in S} a_{ij}$$

BB/FG's formula:

$$Per(A) = \frac{\sum_{\delta} (\prod_{k=1}^n \delta_k) \prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}}{2^{n-1}}$$

$$\delta = \{\delta_1, \delta_2, \dots, \delta_n\} \quad \delta_1 = 1 \text{ and } \delta_i \in \{-1, 1\} \text{ for } 2 \leq i \leq n$$



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# Permanent calculation complexity



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Further complexity improvements by Gray code ordering:

$\mathcal{O}(n^2 \cdot 2^n)$  complexity reduced to  $\mathcal{O}(n \cdot 2^n)$

data recycling in expressions:

$$\sum_{j \in S} a_{ij} \quad \sum_{j=1}^n \delta_j a_{ij}$$

in exchange of scalability reduction?

0	000	0	000
1	001	1	001
2	010	3	011
3	011	2	010
4	100	6	110
5	101	7	111
6	110	5	101
7	111	4	100

## Strawberry Fields

Ryser's formula:

$$\text{Per}(A) = (-1)^n \sum_{S \subseteq \{1,2,\dots,n\}} (-1)^{|S|} \prod_{i=1}^n \sum_{j \in S} a_{ij}$$

## Piquasso

BB/FG's formula:

$$\text{Per}(A) = \frac{\sum_{\delta} (\prod_{k=1}^n \delta_k) \prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}}{2^{n-1}}$$



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# Permanent calculation benchmark



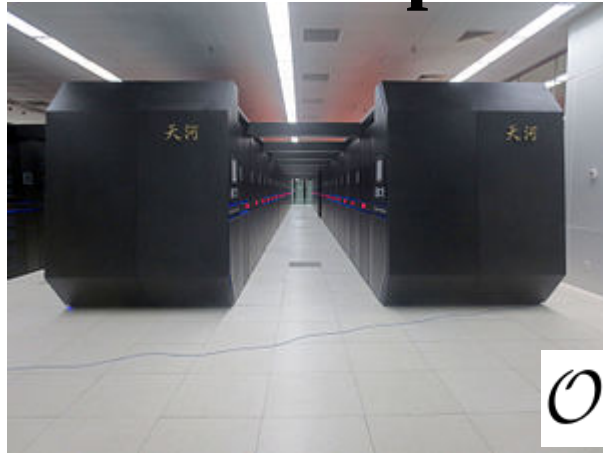
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$$O(n \cdot 2^n)$$

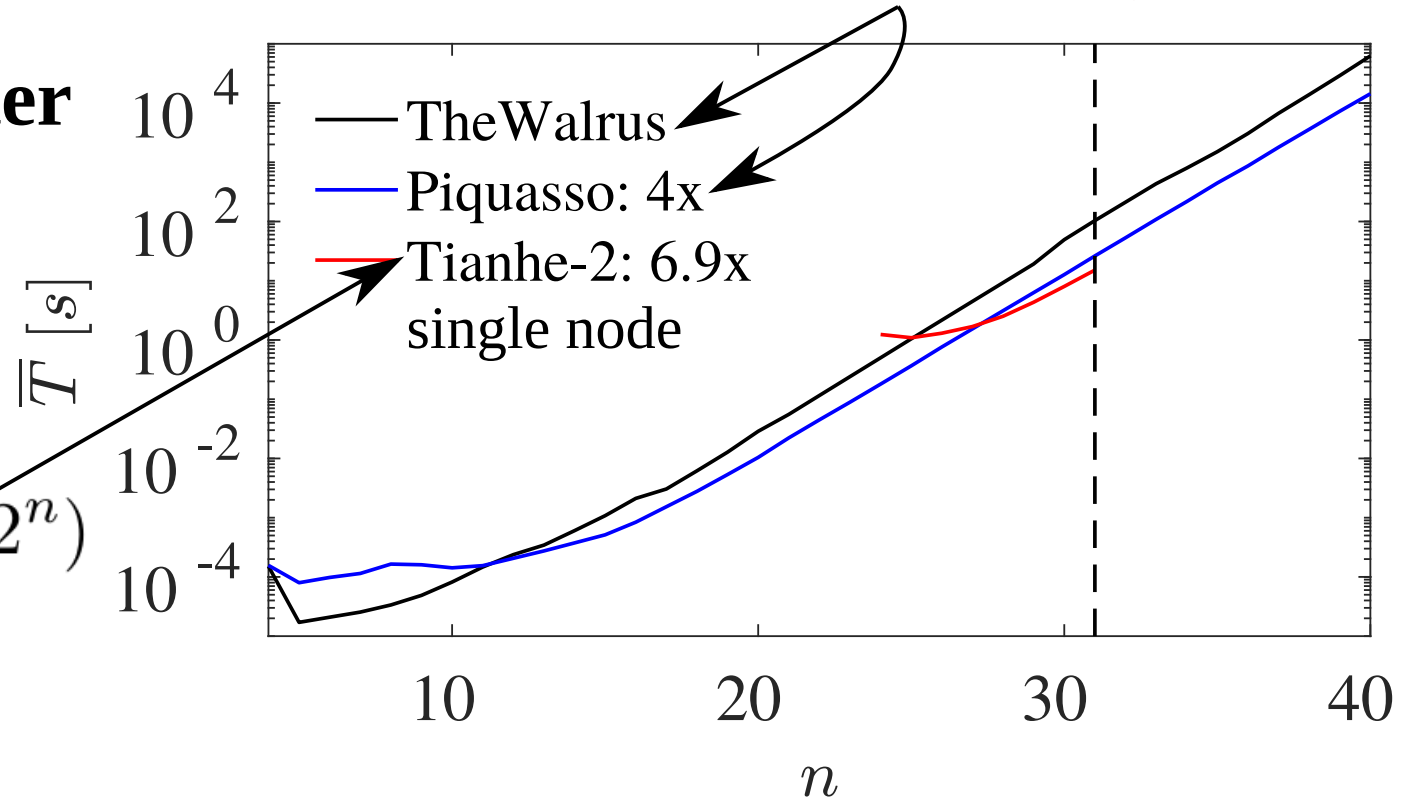
Intel Xeon Gold 6130, 24 threads

## Tianhe-2 supercomputer



$$O(n^2 \cdot 2^n)$$

16 000 computing nodes



**Intel Xeon E5:** dual socket x 12 cores x 2 multithread = 48 threads

**Xeon Phi 31S1P:** 3 PCI cards x 57 cores x 4 multithreads = 684 threads

total: 33.86 petaflops

**limit: permanent of n=50 matrix in 1 hour**

National Science Review, Volume 5, Issue 5, September 2018, Pages 715–720



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# FPGA based data-flow engines (DFE's)



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FPGA: field-programmable gate array

programmable logic gates,  
memory units,  
arithmetic units,  
or other elements



particularly useful for prototyping **application-specific** integrated circuits (ASICs) or processors.

**Quantum Computing inspired ASICs?**

**Coprocessors for real Quantum Hardware?**



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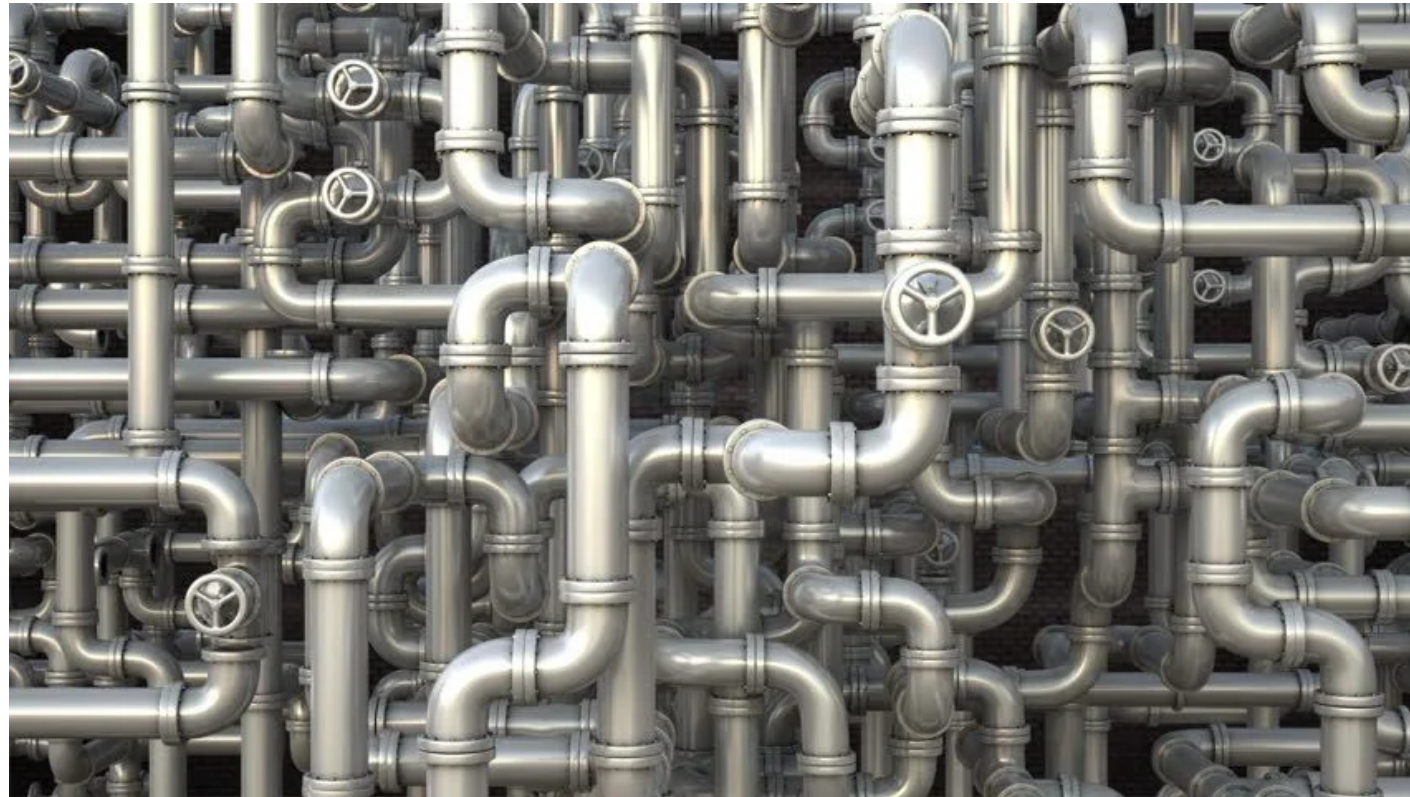
# FPGA based data-flow engines (DFE's)



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Data streams flowing  
through the FPGA chip  
automatized time and  
space constraints



**FPGA hardware + data-flow programming model = DFE**



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supported by Xilinx University Program





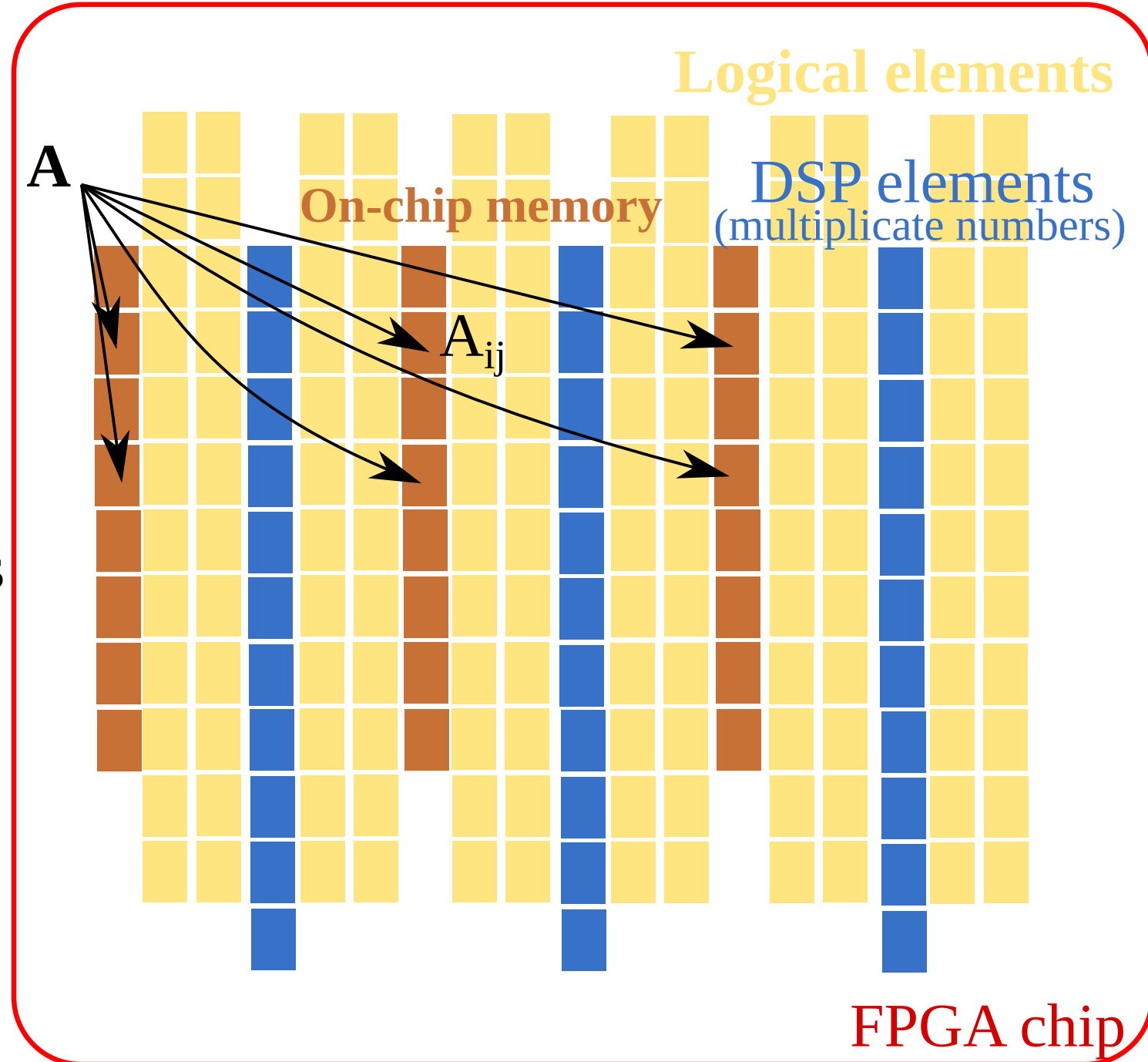
# DFE implementation of permanent calculation

$$\text{Per}(A) = \frac{\sum_{\delta} (\prod_{k=1}^n \delta_k) \prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}}{2^{n-1}}$$

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spread over the  
FPGA chip

**step 1:**  
upload matrix  $A$   
into on-chip memories



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FPGA chip



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# DFE implementation of permanent calculation

$$\text{Per}(A) = \frac{\sum_{\delta} (\prod_{k=1}^n \delta_k) \prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}}{2^{n-1}}$$

$\delta_1 = 1$  and  $\delta_i \in \{-1, 1\}$  for  $2 \leq i \leq n$

counter  $i=1 \dots 2^n \longrightarrow \delta$

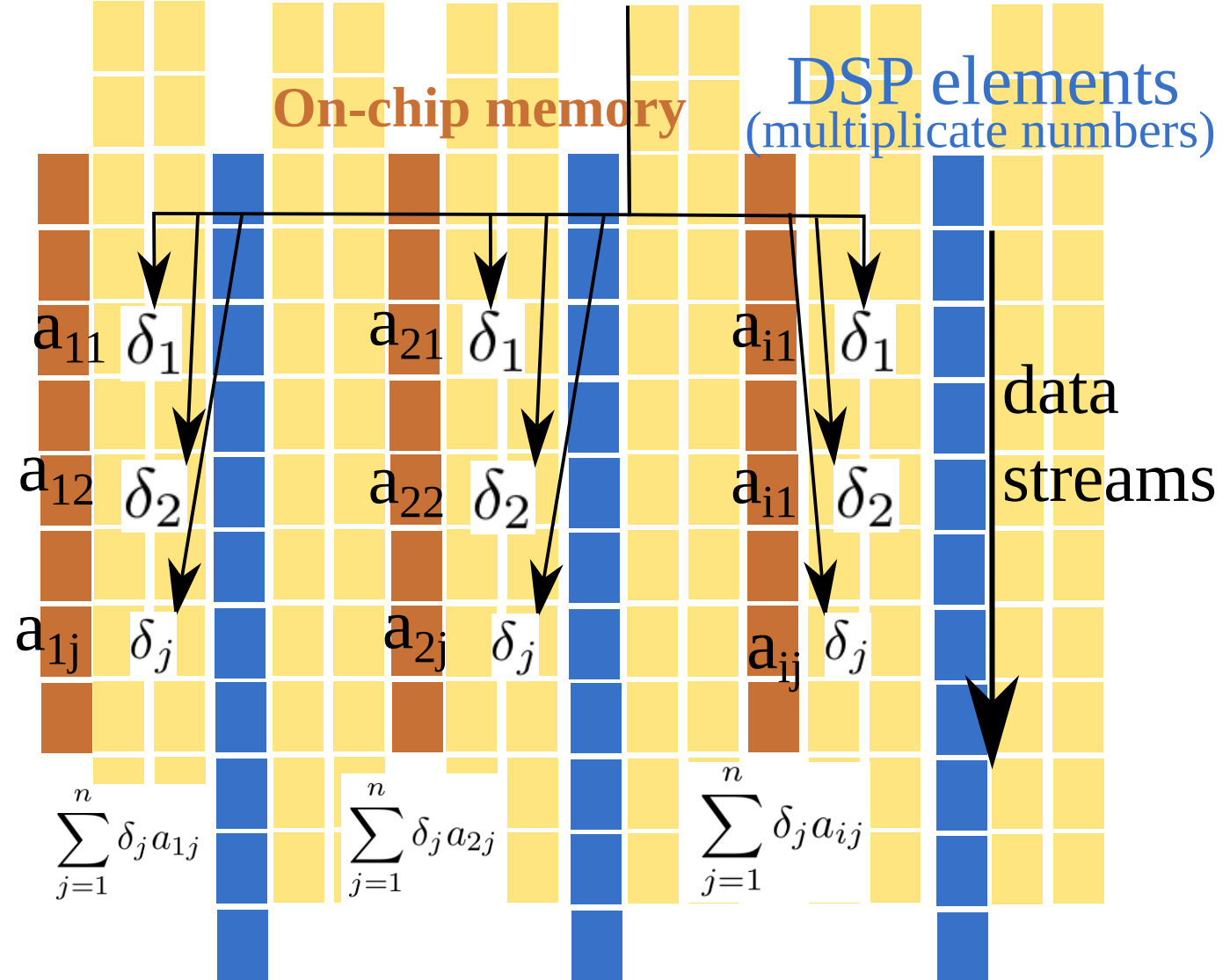
Logical elements

**step 2:**

create counter for  $\delta$

split it into bits:  $\delta_j$

calculate  
for each  $i$   $\sum_{j=1}^n \delta_j a_{ij}$



FPGA chip



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# DFE implementation of permanent calculation

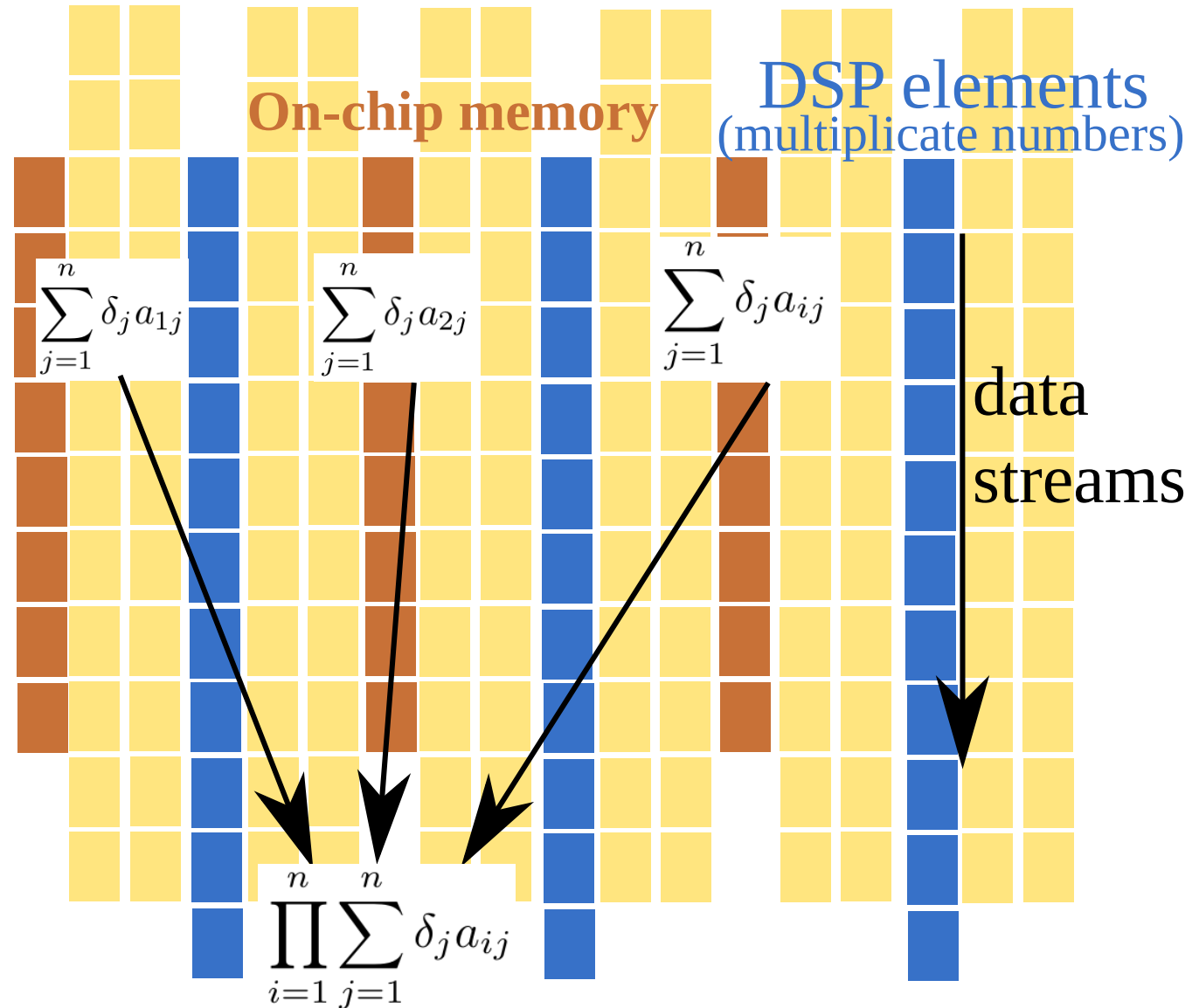
$$\text{Per}(A) = \frac{\sum_{\delta} \left( \prod_{k=1}^n \delta_k \right) \prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}}{2^{n-1}}$$

$\delta_1 = 1$  and  $\delta_i \in \{-1, 1\}$  for  $2 \leq i \leq n$

Logical elements

On-chip memory

DSP elements  
(multiply numbers)



data streams

$$\prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}$$

FPGA chip

**step 3:**

calculate product

$$\prod_{i=1}^n \sum_{j=1}^n \delta_j a_{ij}$$

**step 4:**

sum up the addends  
and return to CPU

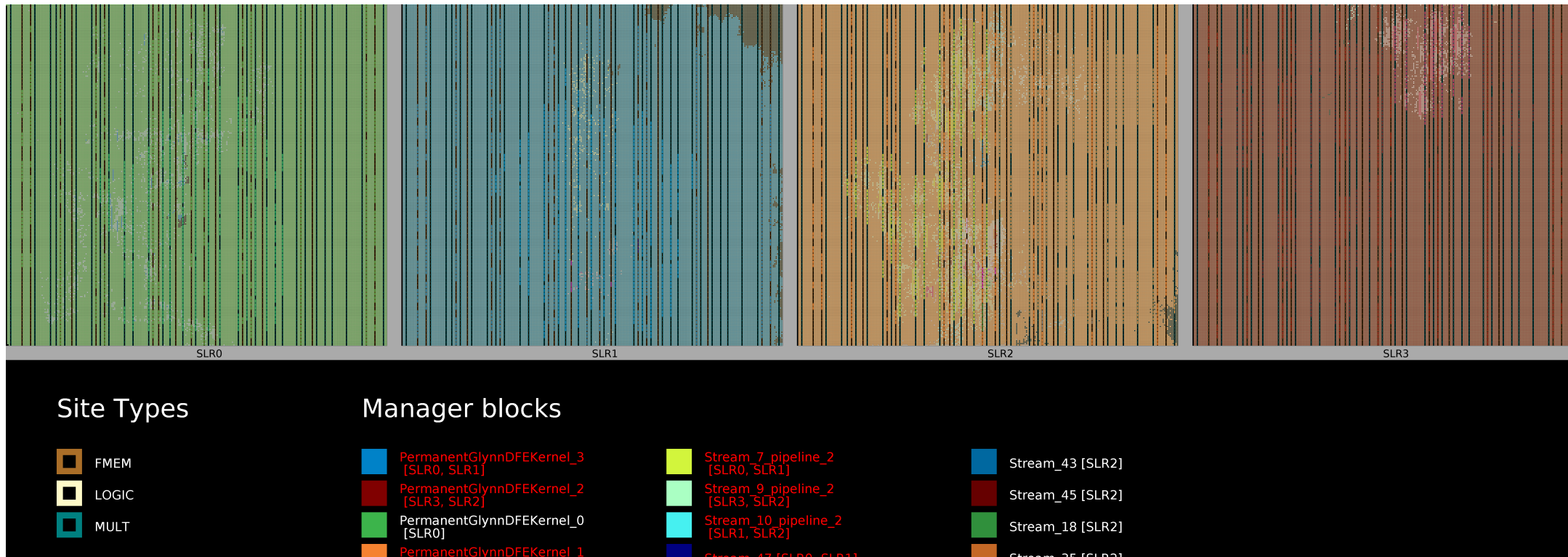


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# DFE implementation of permanent calculation

clock frequency: 180 MHz



Logic utilization: 3750968 / 5184000 (72.36%)

LUTs: 1665146 / 1728000 (96.36%)

Primary FFs: 2085822 / 3456000 (60.35%)

DSP blocks: 4000 / 12288 (32.55%)

Block memory (BRAM18): 3365 / 5376 (62.59%)

Block memory (URAM): 12 / 1280 (0.94%)



# Permanent benchmark again

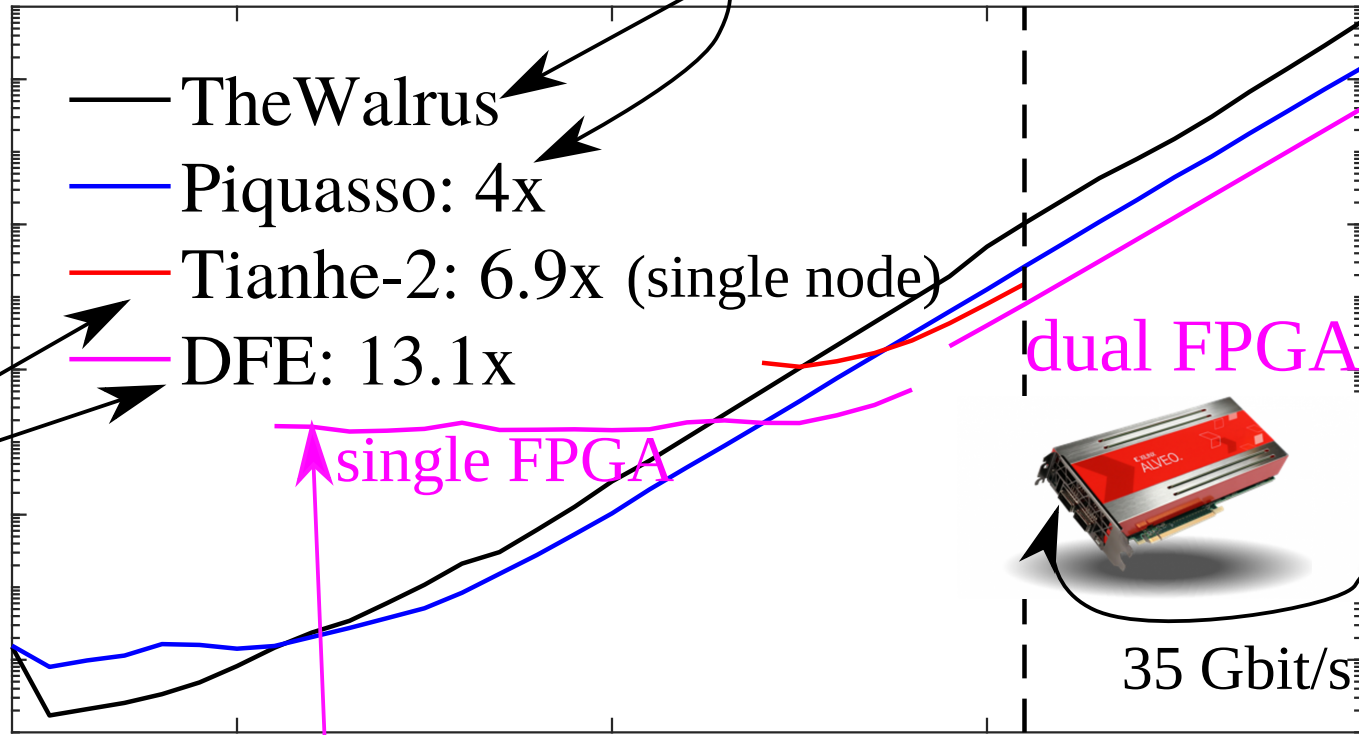


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$\bar{T}$  [s]

$O(n \cdot 2^n)$



$O(n^2 \cdot 2^n)$

single FPGA

dual FPGA

35 Gbit/s transfer

128 bit fixed point  
number representation

uploading the program to FPGA takes time

Time complexity on DFE:  $O(C \cdot 2^n)$

$n^2$  operations are dispersed in space over the FPGA chip



# Numerical error of permanent calculation

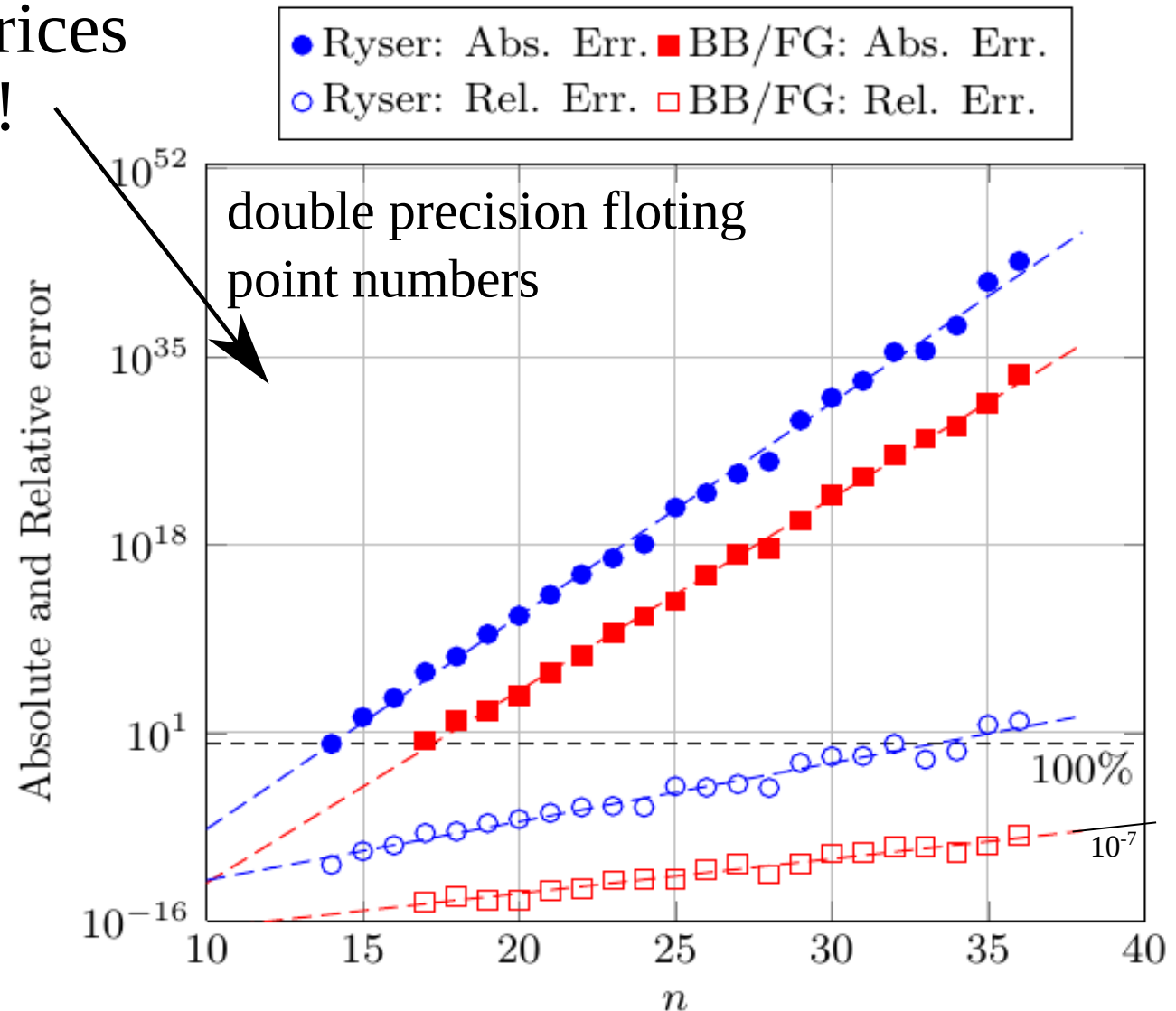
Junjie Wu, Yong Liu, Baida Zhang, Xianmin Jin, Yang Wang, Huiquan Wang, Xuejun Yang, National Science Review, Volume 5, Issue 5, September 2018

$n \times n$  all-ones matrices  
 $\text{Perm}(A) = n!$

If  $A = 40 \times 40$  unitary,

$$a_{ij} \sim \frac{1}{\sqrt{40}}$$

$$\text{Perm}(A) \sim 10^{-12}$$





# Aknowledgement



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