Parallel image reductions over polygon masks

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GPU Day 2024

Motivation

Quality control / selection in manufacturing



A. Types of the defects in metal

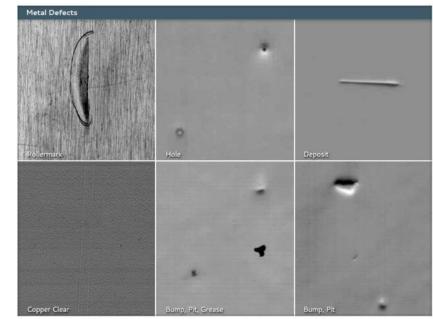


Fig. 1 Different types of surface the defects in metals

Metal Inspection for Surface defect Detection by Image Thresholding D. M. Lohade, P. B. Chopade

Leather defect

Motivation

Quality control / selection in manufacturing

Many goods are produced in wide rolls, sheets, boards, or on conveyor belts

~1-2 meters wide



https://commons.wikimedia.org/wiki/File:Tissue_Paper_Production_Machine.jpg

Desired fault detection size can be as low as 1.0mm - 0.5mm

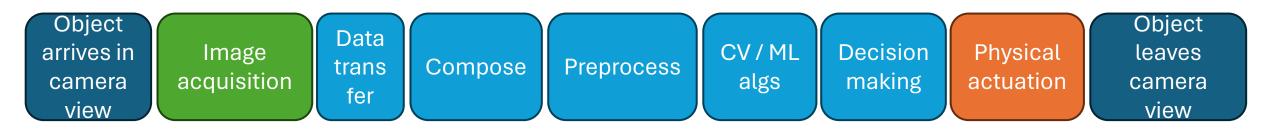
Spatial resolution of up to 10k pixels may be needed

Motivation

Quality control / selection in manufacturing

Movement speed is high, the entire decision making has ~1-10 seconds to decide an accept/reject

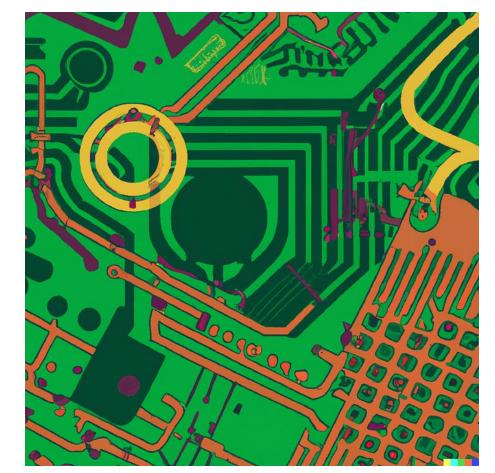
A complex image processing pipeline should finish in this short time



Motivation - polygons

Many applications across industries are working with polygon boundaries, e.g.:

- PCB QC
- Wood QC
- Mining / ore processing
- Civil engineering
- Cartography

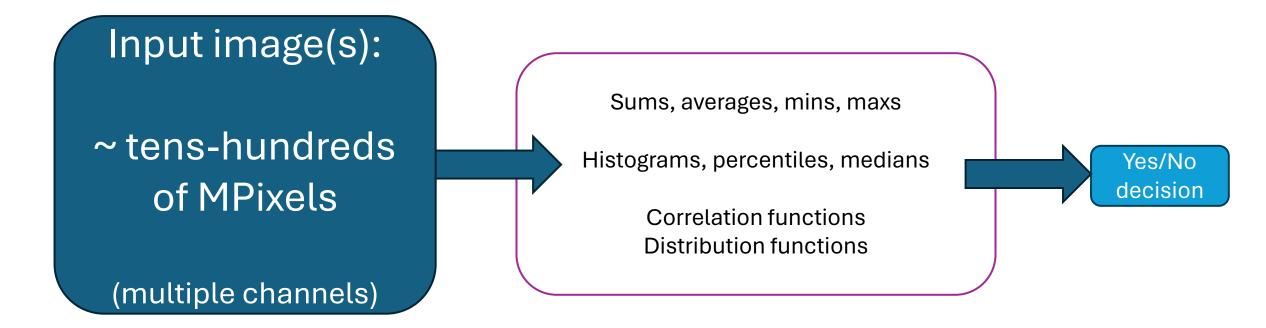


https://medium.com/@nazlicanto/pcb-defect-detection-with-segformer-b56947732914

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Reductions

As with many other image processing tasks, it is all about reductions:



Masked reductions

As with many other image processing tasks, it is all about reductions,

But in many cases, we would like to restrict the area of reductions to some specific region.

Typical masks include: Pixelwise masks, rectangles, or polygons... Sums, averages, mins, maxs

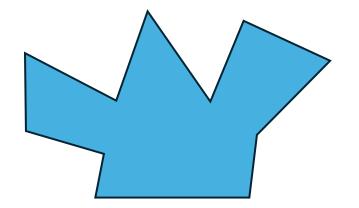
Histograms, percentiles, medians

Correlation functions Distribution functions



Polygons might come from many different sources, we may, or may not make particular assumptions about them.

For the current task, lets assume that we are dealing with <u>simple polygons</u> (no self-intersection, no holes, but can be concave):





Polygons might come from many different sources, we may, or may not make particular assumptions about them.

Let's assume that they are represented as an array of coordinate pairs:

11.73125,-15.850712.015,-15.5694412.51271,-13.4243812.97,-12.78444

Point in polygon problem

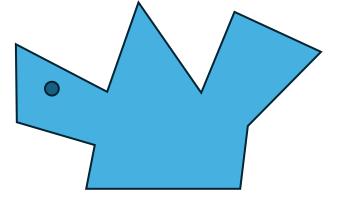
The solution to the masked reduction problem needs to load and aggregate those, and preferably only those pixels that are inside the polygon area.

This is known as the <u>point in polygon</u> problem

There are two widely used algorithms based on:

- Ray casting
- Winding number

Discussion based on: http://profs.ic.uff.br/~anselmo/cursos/CGI/slidesNovos/Inclusion%20of%20a%20Point%20in%20a%20Polygon.pdf



Point in polygon problem

Ray casting:

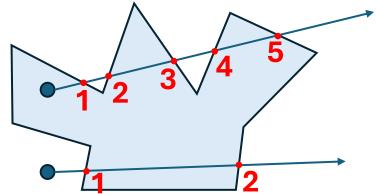
Count the number of intersection points
 with the boundary, if it is odd, the source of the ray is inside

Winding number:

• Sum up the subtended (signed!) angles of the edges If it is non-zero, point is inside

$$\Sigma \theta_i = 0$$

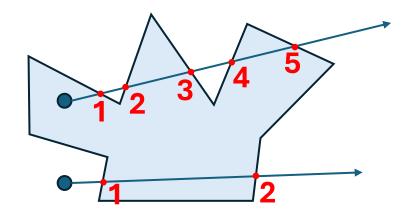
 \wedge



 $\Sigma \theta_i = 2\pi$

Point in polygon problem

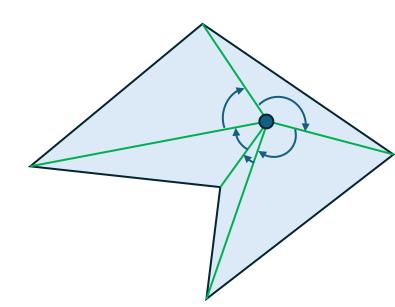
Caveats: Ray casting:



• Intersection computation is prone to many corner cases (like what happens at intersections exactly at the vertices)

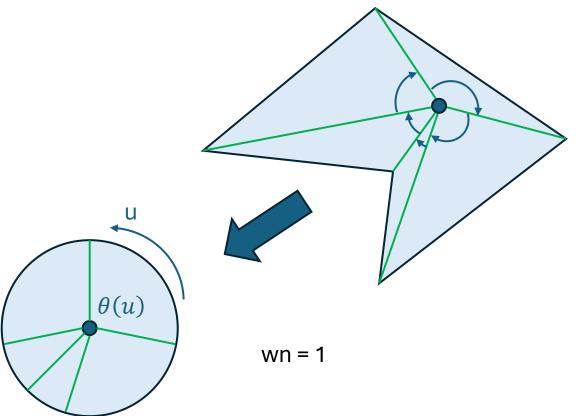
Winding number:

• Less corner case processing is needed



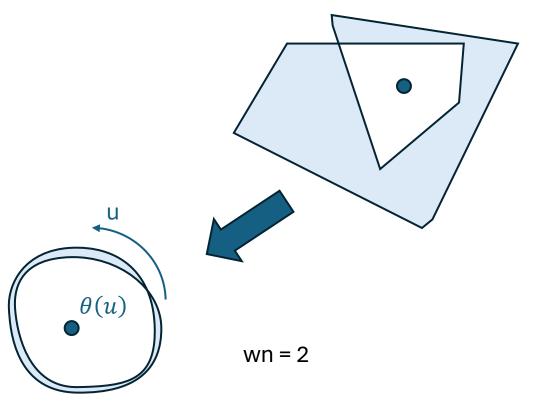
Any closed curve *C* is homeomorphic to the unit circle, so we can use a polar coordinate description to define the winding number:

wn(C) =
$$\frac{1}{2\pi} \oint_C d\theta = \frac{1}{2\pi} \int_0^1 \theta(u) du$$



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For polygons of course, the formulas become more concrete:

 V_2

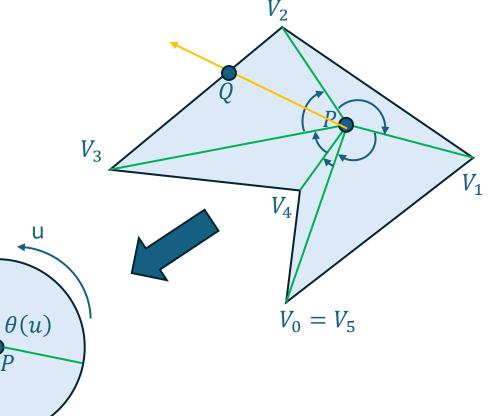
 V_1

$$wn(C) = \frac{1}{2\pi} \oint_{C} d\theta = \frac{1}{2\pi} \int_{0}^{1} \theta(u) du$$

$$= \frac{1}{2\pi} \sum_{i=0}^{n-1} \theta_{i} = \frac{1}{2\pi} \sum_{i=0}^{n-1} a\cos\left(\frac{(V_{i} - P)(V_{i+1} - P)}{|V_{i} - P| |V_{i+1} - P|}\right)$$

In reality, we do not want to compute inverse trigonometric functions, but fortunately we don't need to:

Based on the unit circle correspondence, We only need to track, how many times traversal passes the ray intersection points on the boundary!



In reality, we do not want to compute inverse trigonometric functions, but fortunately we don't need to:

 V_2

 V_{A}

 $V_0 = V_5$

 V_1

 V_3

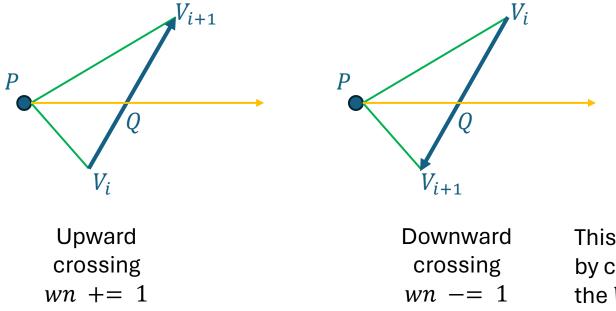
U

 $\theta(u)$

Based on the unit circle correspondence, We only need to track, how many times traversal passes the ray intersection points on the boundary!

This gets awkwardly similar to the ray intersection alg...

But we do not need the intersections, we just need a sign, if an edge is being crossed in one direction, or the other also, we can choose any ray direction, so let's choose a horizontal one:



This relative orientation is obtained by computing the signed area of the $V_iV_{i+1}P$ triangle

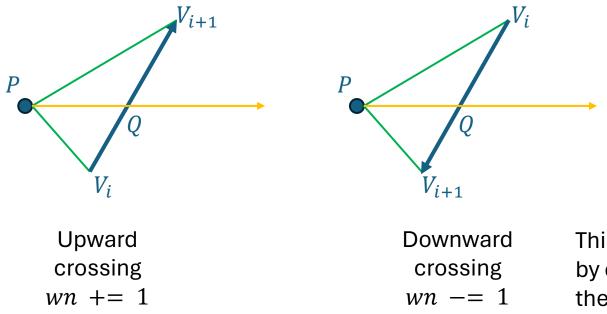
wn(P) = 1

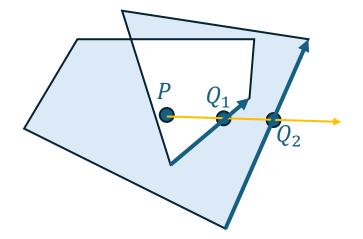
P'

wn(P') = (-1)+(+1) = 0

P

But we do not need the intersections, we just need a sign, if an edge is being crossed in one direction, or the other also, we can choose any ray direction, so let's choose a horizontal one:





This relative orientation is obtained by computing the signed area of the $V_iV_{i+1}P$ triangle

wn(P) = +2

The signed area of a triangle can be obtained by the determinant:

$$A = \frac{1}{2} \begin{vmatrix} 1 & x_0 & y_0 \\ 1 & x_1 & y_1 \\ 1 & x_2 & y_2 \end{vmatrix}$$

i + 1

Ρ

That can also be written:

$$A = \frac{1}{2} \big((x_1 - x_0)(y_2 - y_0) - (x_2 - x_0)(y_1 - y_0) \big)$$

Putting everything together, the following algorithm computes the winding number wn of point P for a polygon given by a list of vertices using a horizontal ray:

```
wn = 0
Loop over edges
Let V0, V1 be the end points of the current edge
if(V1.y <= P.y)
{
    if(V0.y > P.y && area(V1, V0, P) > 0){ wn += 1 }
}
else if(V0.y <= P.y && area(V1, V0, P) < 0){ wn -= 1 }</pre>
```

Problem: We need to test millions of pixels

The naïve algorithm tests every pixel with every edge, this is clearly inefficient

What can we do?

Problem: We need to test millions of pixels

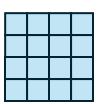
The naïve algorithm tests every pixel with every edge, this is clearly inefficient

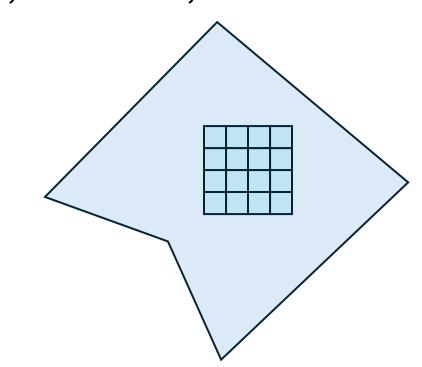
What can we do?

As usual: acceleration structures / early reject

Acceleration structures / early reject:

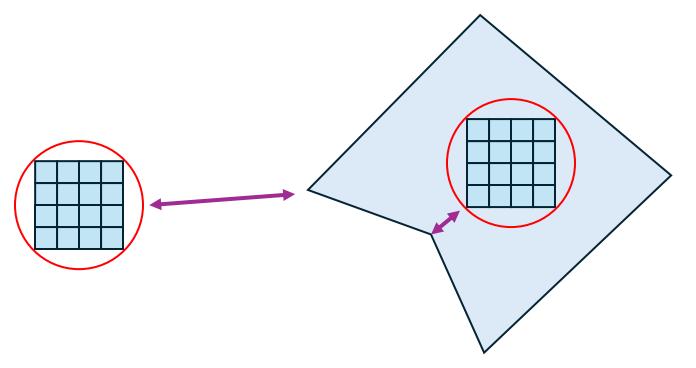
If we have a tile of pixels that is sufficiently far away from the polygon boundaries, it is either entirely inside, or outside, so checking a single point is sufficient





Acceleration structures / early reject:

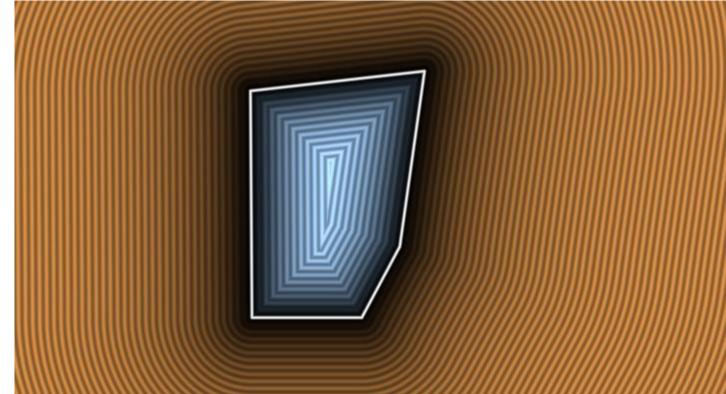
We could compute the minimal distance of the polygon edges to the radius of the circumscribed circle of a square tile:



Signed Distance Fields: an easy to compute function that measures distance from a shape (+) outside

(-) inside

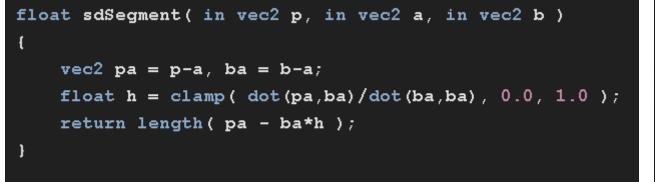
See Inigo Quilez's <u>webpage</u> for formulas and articles

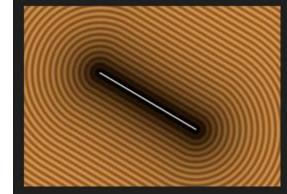


Shadertoy – SDF arbitrary 2D polygon by wagyx

Signed Distance Fields: an easy to compute function that measures distance from a shape

We could take the distance function of segments, and compute the minimum of them





https://iquilezles.org/articles/distfunctions2d

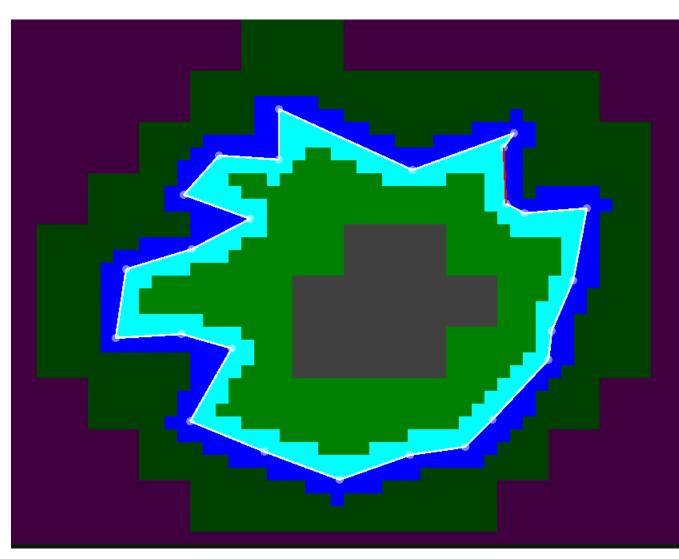
The key insight is the following:

Both the winding number contributions and the line segment SDFs can be computed in parallel for all the edges!

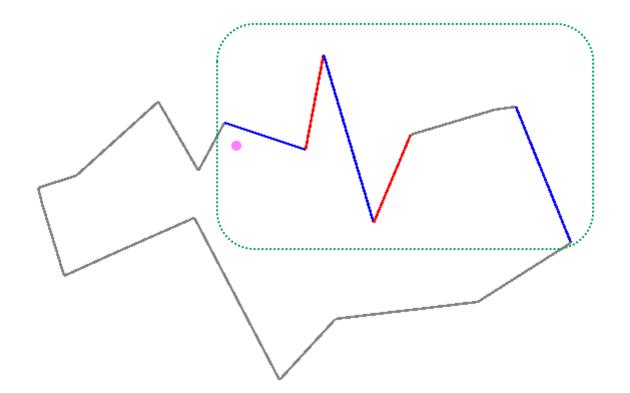
Our threads in a thread block can each pick a segment and compute these simple formulas for the edge's contribution, and we just need to accumulate these partial results. [ray::Vector2 pc = vs[thread];

```
ray::Vector2 pc = vs[thread];
int i2 = thread - 1;
if(thread == 0){ i2 = Nvertices - 1; }
ray::Vector2 pp = vs[i2];
const int wn = wn_line_segment(c, pc, pp);
if(wn!=0){ atomicAdd(&wn_acc, wn); }
d = sdf_line_segment(c, pc, pp);
atomicMin(&dmin_acc, (uint)(1024.0f*d));
```

With large enough tiles, we can reject large areas, but close to the edges we still need to handle all edges for all pixels...

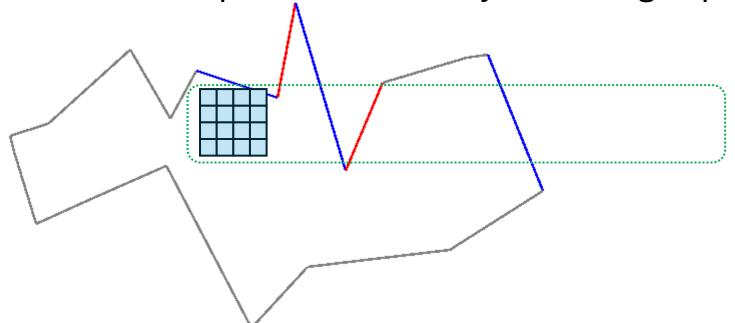


It turns out, that only those edges contribute non-zero to the winding number, that are to the right of the point and whose end's y values are on two sides of the point's y coords.



So, we can filter all the edges (again can be done in parallel) if they are relevant to a particular tile or not:

Usually this narrows down the pixel tests to only a few edges per tile



Results

Let's compare the point in&out rasterization timings (poly: 126 edges):

	RTX 4080 Super 4.15 Mpx	GTX 1650 Mobile 2.07 Mpx	
Naïve:	0.456 ms	2.09 ms	
Tiled:	0.065 ms	0.32 ms	15x-20x
Edge filtered:	0.033 ms	0.22 ms	speed-up
Tiled+Edge filtered:	0.030 ms	0.10 ms 📢	
	30% of ref mem B	3W 68% of ref me	em BW
Reference: vct clear:	0.01 ms	0.068 ms	

Summary

- When porting algorithms to the GPU, first consider the high-level algorithms, and how they can be parallelized efficiently
- Even without micro-optimizations, very large gains are possible over naïve implementations
- Fast image processing codes are in dire need all over the world for various applications

Hungarian GPGPU Community on Discord

- Meeting place for all who are involved in, or would like to learn more of GPGPU technologies
- All fields: Graphics, Compute, Machine Learning, ...
- All APIs: CUDA, OpenCL, Vulkan, OpenGL, ...
- All vendors: AMD, Intel, NVIDIA, ...
- All levels: Students, Teachers, Professionals, Scientists, ...
- All topics: Programming, Bugs, Optimizations, Teaching, Trends, Events, News, ...



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