

Extensions of Digital-Analog Quantum Computation

ReAQCT 2024

Mikel Garcia de Andoin

Tecnalia

Department of Physical Chemistry & EHU Quantum, UPV/EHU

19th of June 2024

Index:

- ① Introduction to DAQC
- ② Digital-analog schedules
- ③ Arbitrary two-body Hamiltonians
- ④ Stability of DAQC for simulations
- ⑤ Conclusions and future directions

Analog quantum computing:

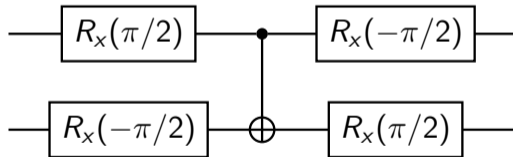
- Continuous evolution under a time-dependent Hamiltonian
- $U_{\text{AQC}} = \mathcal{T} \exp \left(-i \int_0^T dt H(t) \right)$
- Robust (quantum control theory)
- We can only do what the system is designed for

1 - Introduction to DAQC

Gate based/digital quantum computing (DQC):

- Discrete application of unitary operators

- $U_{\text{DQC}} = \prod_i U_i, U_i \in U(2^N)$



- Allows error correction
- Noisy two-qubit gates

“Universal quantum computation can be performed using any entangling interaction and local unitary operations.”

$\{U_{\text{SQG}}, e^{-itH_{\text{TB}}}\}$ is universal

Digital-analog quantum computing (DAQC):

- Combine digital single-qubit gates and analog blocks
- $U_{\text{DAQC}} = \prod_k U_{\text{SQG}} e^{-it_k H_S}$

This allows us to:

- Take advantage of the flexibility of digital quantum computing
- Maintain the robustness against noise of analog quantum computing

2 - Digital-analog schedules

For simplicity, let's use ZZ Hamiltonians

$$H_S = \sum_{i < j}^N h_{i,j} \sigma_i^z \sigma_j^z,$$

$$H_P = \sum_{i < j}^N g_{i,j} \sigma_i^z \sigma_j^z.$$

Objective,

$$e^{-iTH_P} = \prod_k U_k^\dagger e^{-it_k H_S} U_k.$$

2 - Digital-analog schedules

If $H_S \propto H_P$ the schedule is trivial. Else,

$$e^{-iTH_P} = \prod_k U_k^\dagger e^{-it_k H_S} U_k = \prod_k e^{-it_k H_S^{(k)}} = \exp\left(-i \sum_k t_k H_S^{(k)}\right),$$

which is equivalent to

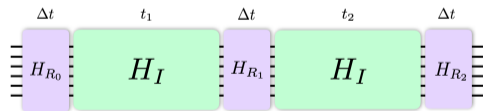
$$M \begin{pmatrix} t_1 \\ t_2 \\ \vdots \\ t_K \end{pmatrix} = T \begin{pmatrix} h_{1,2}/g_{1,2} \\ h_{1,3}/g_{1,3} \\ \vdots \\ h_{N-1,N}/g_{N-1,N} \end{pmatrix}, M \in \{\pm 1\}$$

M has the information about the effective signs of the couplings at each block.

2 - Digital-analog schedules

We assumed a perfect setup $\left\{ \begin{array}{l} \text{Instantaneous SQGs} \\ \text{OR} \\ \text{Ability to turn on and off the interactions on-demand} \end{array} \right.$

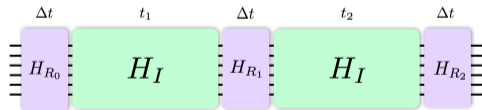
stepwise-DAQC (sDQAC)



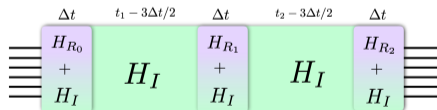
2 - Digital-analog schedules

We assumed a perfect setup $\left\{ \begin{array}{l} \text{Instantaneous SQGs} \\ \text{OR} \\ \text{Ability to turn on and off the interactions on-demand} \end{array} \right.$

stepwise-DAQC (sDQAC)



banged-DAQC (bDAQC)



Banging error:

$$\begin{aligned}U_{\text{Ideal}} = U_{\text{sDAQC}} &= \prod_k e^{-it_k H_S} U_k = \prod_k e^{-it_k H_S} e^{-it_{\text{SQG}} H_{\text{SQG}}^{(k)}} \\ &\approx U_{\text{bDAQC}} = \prod_k e^{-i(t_k - t_{\text{SQG}}) H_S} e^{-it_{\text{SQG}} (H_{\text{SQG}}^{(k)} + H_S)}\end{aligned}$$

For the error to be negligible, $t_{\text{SQG}} \ll \min(t_k)$.
(usually is enough if $10^2 t_{\text{SQG}} \lesssim \min(t_k)$)

3 - Arbitrary two-body Hamiltonians (New!)

3 - Arbitrary two-body Hamiltonians

We wish to work with general two-body Hamiltonians

$$H = \sum_{i,j} \sum_{\mu,\nu \in \{x,y,z\}} h_{i,j}^{\mu,\nu} \sigma_i^\mu \sigma_j^\nu.$$

Before, we had two options to work with arbitrary two-body Hamiltonians:

- Use expensive term-by-term decompositions¹
- Optimize specific cases by hand²

Can we develop a systematic procedure for solving this problem?

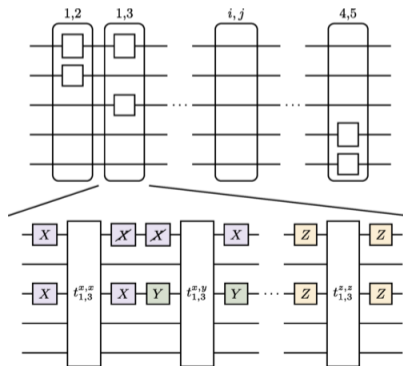
¹A. Parra Rodriguez, P. Lougovski, L. Lamata, E. Solano, and M. Sanz, PRA 101, 022305 (2020)

²T. Gonzalez-Raya, R. Asensio-Perea, A. Martin, et al., PRX Quantum 2, 020328 (2021)

3 - Arbitrary two-body Hamiltonians

Digital-analog quantum computation with arbitrary two-body Hamiltonians

Mikel Garcia-de-Andoin, Álvaro Saiz, Pedro Pérez-Fernández, Lucas Lamata, Izaskun Oregi, and Mikel Sanz
Phys. Rev. Research **6**, 013280 – Published 14 March 2024

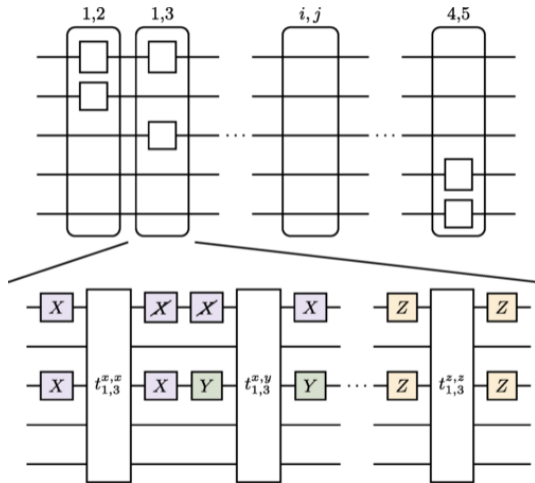


3 - Arbitrary two-body Hamiltonians

By paying a Trotterization error, we treat each possible interaction in the Pauli basis individually.

$$\prod_k e^{-it_k H_S^{(k)}} \approx e^{-i \sum_k t_k H_S^{(k)}}$$

We build a valid DAQC schedule by using combinations of all Pauli gates.

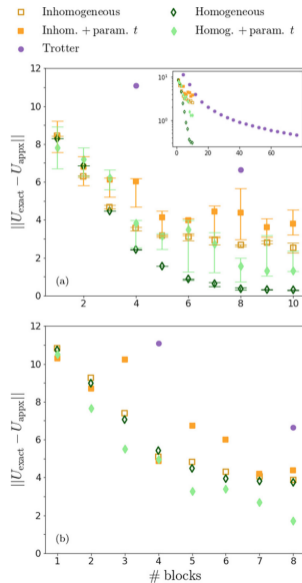


3 - Arbitrary two-body Hamiltonians

This was made assuming that all the interactions we want to simulate $\sigma_i^\mu \sigma_j^\nu$ appear on the source Hamiltonian.

Else,

- “Spray’n’pray”, apply uniformly random single qubit gates.
- Optimize the rotation gates by using a tensor network proxy.



4 - Stability of DAQC (New!)

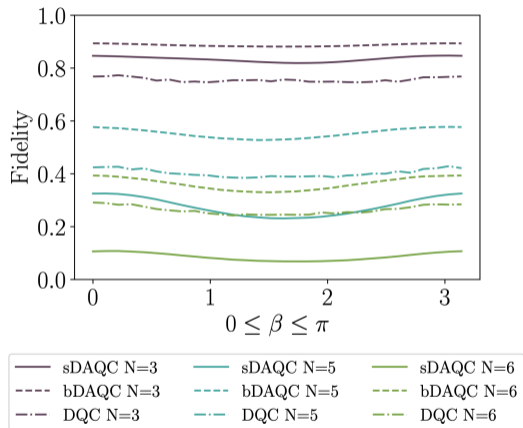
4 - Stability of DAQC

Noise sources:

- SQG errors
- Timing errors
- Bit-flip
- Decoherence
- Dephasing

QFT over a family of states

$$|\Psi(\beta)\rangle = \sin \beta |W_N\rangle + \cos \beta |GHZ_N\rangle$$



We are interested in studying how DAQC scales with the number of qubits.

Informal definition:

A quantum simulation task is stable if the error committed when measuring a local observable depends on the error parameter but doesn't grow with the number of particles.

$$\varepsilon = f(\delta), \quad \varepsilon \neq f(\delta, N)$$

An analogue quantum simulation task on spin systems is stable for local* Hamiltonians.

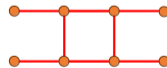
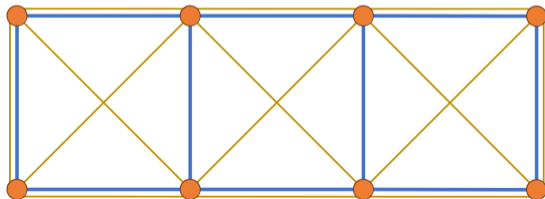
4 - Stability of DAQC

Since SQGs have a very high fidelity, we focus on errors on the source Hamiltonian

$$H_{S,\text{meas}} \longrightarrow H_{S,\text{real}} = H_{S,\text{meas}} + H_{\delta}.$$

This induces certain deviation to the simulated Hamiltonian

$$H_P \longrightarrow H_{P,\text{real}} = H_P + H_{\epsilon}.$$



H_{δ}, H_S, H_P

Since SQGs have a very high fidelity, we focus on errors on the source Hamiltonian, in particular calibration errors

$$H_{S,\text{meas}} \longrightarrow H_{S,\text{real}} = H_{S,\text{meas}} + H_{\delta}.$$

This induces certain deviation to the simulated Hamiltonian

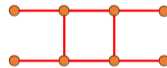
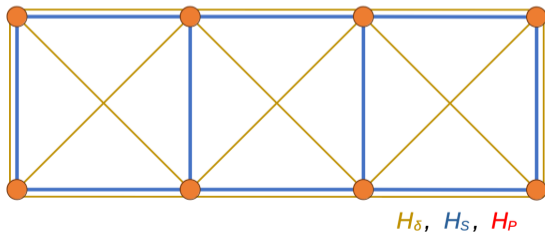
$$H_P \longrightarrow H_{P,\text{real}} = H_P + H_{\epsilon}.$$

4 - Stability of DAQC

If the Hamiltonian is sufficiently local*, the simulation with DAQC is stable.

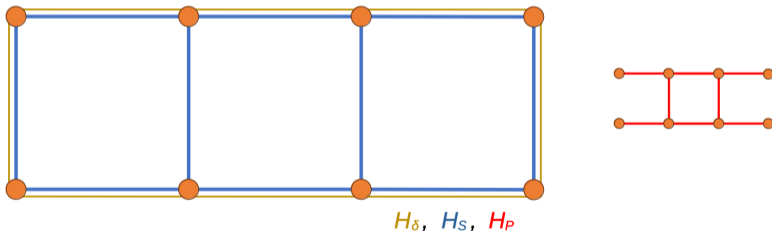
$$\varepsilon \lesssim 4mT\delta\text{deg}(\mathcal{P})\|h_P \otimes h_S\|_\infty + 4m\delta\text{deg}(\mathcal{D}\setminus\mathcal{S})t_A,$$

m number of measured qubits, T simulation time, t_A total time.



We can redesign our DAQC circuit to mitigate errors outside the problem (\sim DD)

$$\varepsilon \lesssim 4mT\delta \deg(\mathcal{P}) \|h_P \otimes h_S\|_\infty.$$



Key points of digital-analog:

- Digital-analog is universal
- Using only SQGs makes it more noise resilient
- Can achieve faster circuit times

Future works:

- Bounds for t_A
- Experimental implementations
- Qudit digital-analog

Affiliations:

tecnal:a

MEMBER OF BASQUE RESEARCH
& TECHNOLOGY ALLIANCE



EHU QC
EHU Quantum Center






nquirephysics.com

Sponsors:



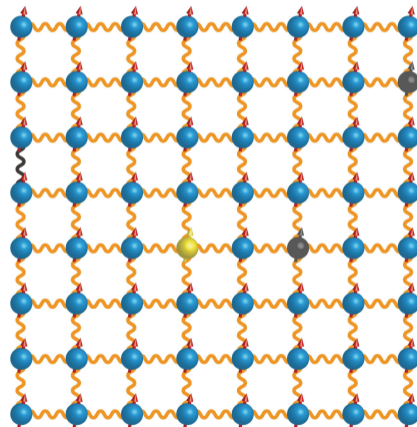
Thank you!

Try DAQC yourself! github.com/NQUIRE-Center/DAQC_simulator 

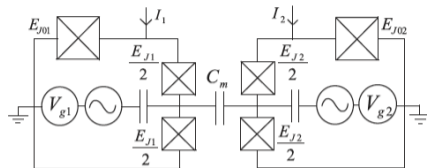
Contact: mikelgda@gmail.com , [@mgarciadeandoin](https://twitter.com/mgarciadeandoin)  

There are already successful experiments employing DAQC

Quantum Neural Networks
experiments on the
Zuchongzhi hardware



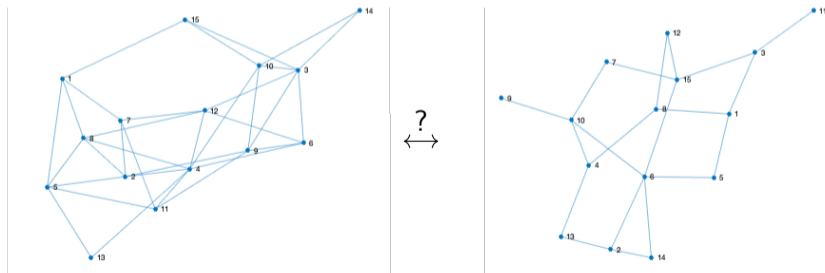
- Superconducting circuits
- Trapped ions
- Rydberg atoms
- Liquid state NMR setups
- ...



$$\begin{aligned}
 H = & \frac{\omega_1}{2} \sigma_1^z + \Omega_1 \cos(\omega_1^{\text{rf}} t + \phi_1) \sigma_1^x \\
 & + \frac{\omega_2}{2} \sigma_2^z + \Omega_2 \cos(\omega_2^{\text{rf}} t + \phi_2) \sigma_2^x \\
 & + \frac{\omega_{xx}}{2} \sigma_1^x \sigma_2^x. \\
 \rightarrow \dots \rightarrow & H_{QF}^{\text{eff}} \sim \omega_{xx} \sigma_1^x \sigma_2^x
 \end{aligned}$$

Simulation of dynamics of arbitrary Hamiltonians on an arbitrary system:

- If the connectivity graph of H_P doesn't match the one in H_S we need to implement SWAP strategies (really hard).



[A Scheme to Implement a Universal Two-Qubit Quantum Circuit using Cross-Resonance Hamiltonian.pdf](#)

[Benchmarking Digital-Analog Quantum Computation.pdf](#)

[Digital-analog co-design of the Harrow-Hassidim-Lloyd algorithm.pdf](#)

[Digital-Analog Counterdiabatic Quantum Optimization with Trapped Ions.pdf](#)

[Digital-analog for the Fourier quantum transform.pdf](#)

[Digital-Analog Quantum Computation.pdf](#)

[Digital-analog quantum convolutional neural networks for image classification.pdf](#)

[Digital-analog quantum learning on Rydberg atom arrays.pdf](#)

[Digital-Analog Quantum Simulations Using the Cross-Resonance.pdf](#)

[Enhanced connectivity of quantum hardware with digital-analog control.pdf](#)

[Enhancing Quantum Annealing in Digital-Analog Quantum Computing.pdf](#)

[Fixed Depth Hamiltonian Simulation via Cartan Decomposition.pdf](#)

[Hybrid DA simulation of many-body dynamics with SC qubits.pdf](#)

[Noise in Digital and Digital-Analog Quantum Computation.pdf](#)

[Qadence- a differentiable interface for digital-analog programs.pdf](#)

[Quantum Computing with an Always-On Heisenberg Interaction.pdf](#)

[Realization of programmable Ising models in a trapped-ion quantum simulator.pdf](#)

[Stochastic error cancellation in analog quantum simulation.pdf](#)

2 - Digital-analog schedules (example)

We want to simulate a ZZ-Ising Hamiltonian for a time $T=1$, with another ZZ-Ising Hamiltonian.

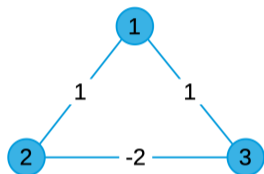


Figure: Source Hamiltonian (H_S)

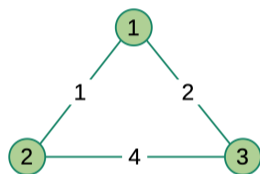


Figure: Problem Hamiltonian (H_P)

$$H_S = 1\sigma_1^z\sigma_2^z + 2\sigma_1^z\sigma_3^z - 2\sigma_2^z\sigma_3^z,$$

$$H_P = 1\sigma_1^z\sigma_2^z + 2\sigma_1^z\sigma_3^z + 4\sigma_2^z\sigma_3^z,$$

2 - Digital-analog schedules (example)

By sandwiching an X gate to a single qubit, we can change the effective sign of the couplings connected to it.

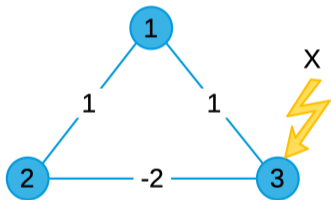


Figure: Sandwich block with X gate on qubit 3

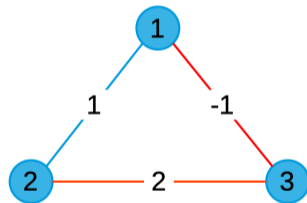


Figure: Effective Hamiltonian ($H_S^{(k)}$)

2 - Digital-analog schedules (example)

Run all possible combinations of SQGs to obtain

$$M = \begin{pmatrix} 1 & -1 & -1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \end{pmatrix}$$

Find the times for the analog blocks.

$$\begin{pmatrix} 1 & -1 & -1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \end{pmatrix} \begin{pmatrix} t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = \begin{pmatrix} h_{1,2}/g_{1,2} \\ h_{1,3}/g_{1,3} \\ h_{2,3}/g_{2,3} \end{pmatrix} = \begin{pmatrix} 1/1 = 1 \\ 2/2 = 1 \\ 4/-2 = -2 \end{pmatrix}$$

We want to minimize the total time of the circuit, and the times to be positive

$$\min_{\vec{t}} \sum_k t_k \quad \oplus \quad M \cdot \vec{t} = Th/g \quad \& \quad t_k \geq 0 \quad \forall k$$

2 - Digital-analog schedules (example)

The solution to this problem is

$$t_1 = 1, t_2 = 0, t_3 = 1.5, t_4 = 1.5.$$

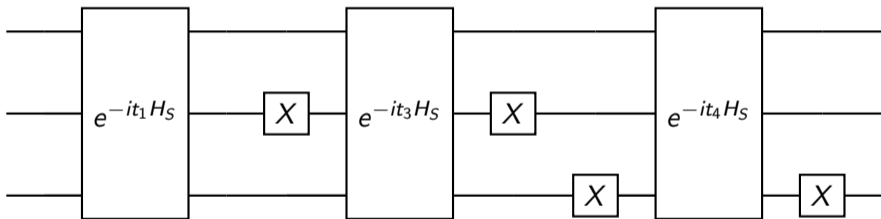
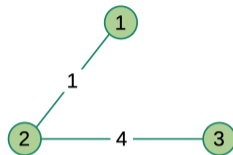
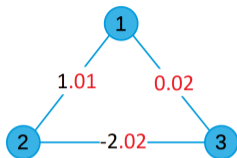


Figure: DAQC circuit for simulating H_P for $T=1$ with H_S .

4 - Stability of DAQC



Time optimal: eliminate non existent couplings

$$\begin{pmatrix} 1 & -1 & -1 & 1 \\ 1 & -1 & 1 & -1 \\ \cancel{1} & \cancel{1} & \cancel{-1} & \cancel{-1} \end{pmatrix} \begin{pmatrix} t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = T \begin{pmatrix} h_{1,2}/g_{1,2} \\ h_{1,3}/g_{1,3} \\ \cancel{0/0} \end{pmatrix}$$

Error mitigated: set non existent couplings to 0 (cost extra circuit time)

$$\begin{pmatrix} 1 & -1 & -1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \end{pmatrix} \begin{pmatrix} t_1 \\ t_2 \\ t_3 \\ t_4 \end{pmatrix} = T \begin{pmatrix} h_{1,2}/g_{1,2} \\ h_{1,3}/g_{1,3} \\ 0/0 = 0 \end{pmatrix}$$