

COMPUTING CLASSICAL PARTITION FUNCTIONS: FROM ONSAGER AND KAUFMAN TO QUANTUM ALGORITHMS

ReAQCT '24, June 19–20, 2024 | Roberto Gargiulo, Matteo Rizzi, Robert Zeier | Quantum Control (PGI-8), FZJ



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The Context

Optimization, Partition Functions and Quantum Computers



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- Binary Optimization (QUBO/MaxCut) is hard
- Many Classical Approaches:
 - Heuristic Exact Solvers (e.g. Gurobi, QuBowl)
 - Hardware Solvers (e.g. Ising Machines)
 - Monte Carlo (e.g. Simulated Annealing)
 - More...
- State-of-the-art: $\sim 10^4$ vertices (sparse graphs)





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- Quantum Approaches:
 - Quantum Annealing (D-Wave)
 - Variational Algorithms (QAOA, VQE)
 - Imaginary Time Evolution (VITE, QITE)
 - And many more







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 - And many more
- Classical-to-Quantum Optimization Problem:

$$\begin{split} \{ \boldsymbol{s}_{\boldsymbol{v}} \} &\in \{ \boldsymbol{0}, \boldsymbol{1} \}^{M} \to |\psi\rangle \in (\mathbb{C}^{2})^{\otimes M} \\ \min_{\{ \boldsymbol{s}_{\boldsymbol{v}} \}} \boldsymbol{H}(\{ \boldsymbol{s}_{\boldsymbol{v}} \}) \to \min_{\psi} \langle \psi | \boldsymbol{H} | \psi \rangle \end{split}$$

• State-of-the-art: $\sim 10^2 - 10^3$ qubits $\lesssim 10^4$







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- State-of-the-art: $\sim 10^2-10^3~\text{qubits} \lesssim 10^4$
- NISQ: Limited Memory (and Circuit Depth). Alternative road?







BEYOND OPTIMIZATION

- Physics-inspired approach:
 Finite temperature Ising models
- Still hard at finite temperature
- Classical Approaches:
 - Monte Carlo Methods
 - Tensor Networks
 - Hardware Solvers (Janus II)
 - And many more
- Alternative Classical-to-Quantum Mapping: Solution of 2D Ising Model







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The Prelude

The Classical 2D Ising Model



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THE PRELUDE - THE CLASSICAL 2D ISING MODEL

• **Classical** spins on a 2D lattice with $s_{v,k} = \pm 1$ and energy

$$E(\{s_{v,k}\}) = -\sum_{k=1}^{N} \sum_{\nu=1}^{M} (Js_{v,k}s_{\nu+1,k} + \Gamma s_{\nu,k}s_{\nu,k+1}), J, \Gamma > 0$$

- Equilibrium Properties, $p(\{s_{v,k}\}) \propto e^{-\beta E(\{s_{v,k}\})}$:
 - **1** Partition Function: $\mathcal{Z} = \sum_{\{s_{v,k}\}} e^{-\beta E(\{s_{v,k}\})}$
 - **2** Correlation Functions: $\langle s_{u,k}s_{u',k'} \rangle = \frac{1}{\mathcal{Z}} \sum_{\{s_{v,k}\}} s_{u,k}s_{u',k'} e^{-\beta E(\{s_{v,k}\})}$





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Transfer Matrix Solution (Onsager 1944, Kaufman 1949):

- **1** Partition Function: $\mathcal{Z} = c \operatorname{Tr}\left(\prod_{k=1}^{N} V\right), V = e^{\beta J \sum_{\nu=1}^{M} Z_{\nu} Z_{\nu+1}} e^{\gamma \sum_{\nu=1}^{M} X_{\nu}}$
- 2 Correlation Functions: $\langle s_{u,k} \hat{s}_{u',k'} \rangle = \frac{1}{Z} \operatorname{Tr}(V \cdots Z_u V \cdots Z_{u'} V \cdots V)$
- (Floquet) Imaginary Time Evolution of the Quantum Chain
- Only *M* qubits for $M \times N$ spins!





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- **2** Correlation Functions: $\langle s_{u,k} s_{u',k'} \rangle = \frac{1}{Z} \operatorname{Tr}(V \cdots Z_u V \cdots Z_{u'} V \cdots V)$
- Diagonalization via Representation Theory of Lie Groups/Algebras:

(aka Jordan-Wigner and Free Fermions) $\rho \colon \mathrm{SO}(2M,\mathbb{C})^{\times 2} \mapsto \mathcal{M}(2^M,\mathbb{C})$ $T = P \cdot \begin{pmatrix} t_1 \\ & \ddots \\ & & t_{4M} \end{pmatrix} \cdot P^{-1} \mapsto V = \rho(T) = \rho(P) \begin{pmatrix} v_1 \\ & \ddots \\ & & v_{2^M} \end{pmatrix} \rho(P^{-1}) \Rightarrow \mathbb{Z}$



MOTIVATION AND OBJECTIVES

Motivation

- 1 Partition Functions as a **generalization** of optimization
- 2 Qubit-efficient Classical-to-Quantum Mapping for 2D Ising Model¹

Objectives

- More general models? Kaufman-type/Lie-theoretic solution?
- Can we (efficiently) implement it on a quantum computer?

¹See also: Arad 2010, De las Cuevas 2011, Iblisdir 2014, Matsuo 2014



• **Classical** spins $s_{v,k} = \pm 1$ on *N* layers of a graph \mathbb{G} :

$$E(\{s_{v,k}\}) = -\sum_{k=1}^{N} \left(\sum_{(u,v)\in E(\mathbb{G})} J_{uv}^{k} s_{u,k} s_{v,k} + \sum_{v\in V(\mathbb{G})} H_{v}^{k} s_{v,k} + \sum_{v\in V(\mathbb{G})} \Gamma_{v}^{k} s_{v,k} s_{v,k+1} \right)$$

Equilibrium Properties via Transfer Matrix (here Γ^k_ν > 0):
 Partition Function (Periodic Boundary):

$$\mathcal{Z} = c \operatorname{Tr}\left(\prod_{k=1}^{n} V_{k}\right), \quad V_{k} = e^{\beta \sum_{(u,v)} J_{uv}^{k} Z_{u} Z_{v}} e^{\beta \sum_{v} H_{v}^{k} Z_{v}} e^{\sum_{v} \gamma_{v}^{k} X_{v}}$$

2 Correlation Functions (Periodic Boundary):

$$\langle s_{u,k}s_{u',k'}\rangle = \frac{1}{\mathcal{Z}}\operatorname{Tr}(V_1\cdots Z_u V_k\cdots Z_{u'} V_{k'}\cdots V_N)$$

Dimensional Reduction: $M \times N$ classical spins to M qubits.

• N^d spins on *d*-dim hypercube $\rightarrow N^{d-1}$ qubits on *d* - 1-dim hypercube



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• (NP-)Hard beyond standard 2D:

1 Additional terms: Square 2D, with fields $H_v^k \neq 0$;

2 Increased connectivity: Cubic 3D, no fields $H_{\nu}^{k} = 0$.

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- Kaufman-type/Lie-theoretic solution? Exponential time/memory:
 - **1** For any graph \mathbb{G} , with fields, dim $G_{\text{Lie}} = O(4^M)$

2 For any non-1D graph \mathbb{G} , no fields, dim $G_{\text{Lie}} = O(4^M)^{(1)}$

⁽¹⁾Real-imaginary correspondence based on the work of Kazi/Larocca/Farinati/Coles/Cerezo/Zeier (Unpublished)



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Probabilistic Approach: Block Encodings²



Introduce non-unitarity via ancillas and (weak) measurements

²See also: Martyn 2021, Zhu 2023, Arad 2010.

- Transfer Matrix V_k is **not unitary**! No direct implementation on quantum computers
- Probabilistic Approach: Block Encodings



- (1) Many ancillas, end-of-circuit measurements
- (2) Few ancillas, mid-circuit measurements with reset

• Transfer Matrix V_k is **not unitary**! No direct implementation on quantum computers

Deterministic Approach: Unitary Approximation (Open Boundary)

1 Variational Imaginary Time Evolution (McArdle, 2019) as subroutine:



Many-ancilla Block-Encoding

(Periodic Boundary: $\rho \propto 1$, Open Boundary: $\rho = (|+\rangle \langle +|)^{\otimes M}$)

1 Partition Function via **Hadamard Test** with $\Pi_0^C + \Pi_1^C U_V, V = \prod_k V_k$:

$$\mathcal{Z} = c \operatorname{Tr}(\rho \mathbf{V}) = c'(\langle X_C \rangle - \mathbf{i} \langle Y_C \rangle)$$



2 Expectation values via Hadamard Test with $\Pi_0^C + \Pi_1^C U_{V'}, V' = V_1 \cdots Z_u V_k \cdots Z_{u'} V_{k'} \cdots V_N$:

$$\langle s_{u,k} s_{u',k'} \rangle = \frac{\operatorname{Tr}(\rho V')}{\operatorname{Tr}(\rho V)} = \frac{\langle X_C \rangle' - i \langle Y_C \rangle'}{\langle X_C \rangle - i \langle Y_C \rangle}$$



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- **Few-Ancilla** Block-Encoding (Stochastic Circuit) (Periodic Boundary: $\rho \propto 1$, Open Boundary: $\rho = (|+\rangle \langle +|)^{\otimes M}$)
 - **1** Partition Function via **Hadamard Test** with $\Pi_0^C + \Pi_1^C V$, $V = \prod_k V_k$:

$$\mathcal{Z} = c \operatorname{Tr}(\rho V) = c' \frac{\langle X_C \rangle - i \langle Y_C \rangle}{1 + \langle Z_C \rangle}$$



$$\langle \boldsymbol{s}_{u,k} \boldsymbol{s}_{u',k'} \rangle = \frac{\operatorname{Tr}(\rho \,\boldsymbol{V}')}{\operatorname{Tr}(\rho \,\boldsymbol{V})} = \frac{\langle \boldsymbol{X}_{\mathcal{C}} \rangle' - \boldsymbol{i} \, \langle \boldsymbol{Y}_{\mathcal{C}} \rangle'}{\langle \boldsymbol{X}_{\mathcal{C}} \rangle - \boldsymbol{i} \, \langle \boldsymbol{Y}_{\mathcal{C}} \rangle} \frac{1 + \langle \boldsymbol{Z}_{\mathcal{C}} \rangle}{1 + \langle \boldsymbol{Z}_{\mathcal{C}} \rangle'}$$





- Unitary Approximation (Open Boundary: $\rho = (|+\rangle \langle +|)^{\otimes M}$)
 - 1 Partition Function via **Hadamard Test** with $U(\theta^*) \approx \Pi_0^C + \Pi_1^C V, V = \prod_k V_k$:

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SUMMARY AND OUTLOOK

Transfer Matrix mapping for equilibrium classical systems on quantum systems:

- Dimensional Reduction
- Works beyond standard 2D.
- No Kaufman-type solution beyond standard 2D
- 2 Quantum Computer implementations suitable for NISQ:
 - Block encodings: Polynomial depth, Variable number of ancillas
 - Unitary Approximation: No ancillas, Model-dependent depth
 - Approximation Scale?
 - General inter-layer interactions?
 - Role of Symmetries? Lie-theoretic properties?



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Thanks!



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LIE THEORY IN A NUTSHELL

The imaginary time circuit V = ∏_k e<sup>∑_(u,v) α^k_{u,v}Z_vZ_v e^{∑_v β^k_vZ_v} e^{∑_v γ^k_vX_v} belongs to a (real) Lie group G ⊆ GL(2^M, C)
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- The infinitesimal generators $\mathcal{G} = \{Z_u Z_v\} \cup \{X_v\} \cup \{Z_v\}$ define the Lie algebra \mathfrak{g} :

 $\mathfrak{g} = \mathsf{Span}_{\mathbb{R}} \{ \mathcal{G} \cup \{ [\mathbf{Z}_{\boldsymbol{U}} \mathbf{Z}_{\boldsymbol{v}}, \mathbf{X}_{\boldsymbol{v}'}] \} \cup \{ [\mathbf{X}_{\boldsymbol{v}}, \mathbf{Z}_{\boldsymbol{v}'}] \} \cup \cdots \}$

Study G (a group) by looking at g (an algebra), e.g.: dimension, invariant subspaces, decompositions, reachability.



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FROM IMAGINARY TO REAL: WICK ROTATION

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- Given an **involution** θ : $\mathfrak{g} \mapsto \mathfrak{g}$, $\theta^2 = \mathbb{1}_{\mathfrak{g}}$ on a (semisimple) Lie algebra:

$$\mathfrak{g} = \mathfrak{l} \oplus \mathfrak{p}, \ \theta(\mathfrak{l}) = +\mathfrak{l}, \ \theta(\mathfrak{p}) = -\mathfrak{p}$$

one can go from a Lie algebra to another (with $\mathfrak{g}_{\mathbb{C}}=\tilde{\mathfrak{g}}_{\mathbb{C}})$:

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- **Multiple choices** of *θ* (Wick rotation) are possible!
- Special cases: one compact and one split real form.



LIE ALGEBRA CLASSIFICATION

Conjecture

The "thermal" Lie algebra $\mathfrak{g} = \langle \{Z_u Z_v\}_{(u,v) \in E(G)}, \{X_v\}_{v \in V(G)} \rangle_{\text{Lie}}$ for classical layers of arbitrary graphs *G* (no fields) is the **split real form** of the Multi-Angle QAOA ansatz.

(Imaginary Time) Disordered Ising model on:

- 2D Rectangular lattice with periodic boundary conditions: $g = so(M, M) \oplus so(M, M)$
- 3D (even) Cubic lattice Ising model: $\mathfrak{g} = \mathfrak{sp}(2^{M-1}, \mathbb{R}) \oplus \mathfrak{sp}(2^{M-1}, \mathbb{R})$

(**Real Time**) Multi-Angle QAOA on⁽¹⁾:

- Cycle Graph: g = so(2M) ⊕ so(2M)
- (Even) Bipartite Graph: $\mathfrak{g} = \mathfrak{sp}(2^{M-1}) \oplus \mathfrak{sp}(2^{M-1})$

⁽¹⁾(Unpublished) Work by Kazi/Larocca/Farinati/Coles/Cerezo/Zeier, 2024

LIE ALGEBRA CLASSIFICATION

Result #1

The "thermal" Lie algebra $\mathfrak{g} = \langle \{Z_u Z_v\}_{(u,v) \in E(G)}, \{X_v, Z_v\}_{v \in V(G)} \rangle_{\text{Lie}}$ for classical layers of any connected graph *G* (with fields) is $\mathfrak{sl}(2^M, \mathbb{R}), M = |V|$.

Result #2

The "thermal" Lie algebra $\mathfrak{g} = \langle \{Z_u Z_v\}_{(u,v) \in E(G)}, \{X_v\}_{v \in V(G)} \rangle_{\text{Lie}}$ for classical layers of arbitrary graphs *G* (no fields) has the **same dimension** as that of the Multi-Angle QAOA ansatz.