Scalable volumetric benchmarks based on Clifford and free-fermion operations

Orsolya Kálmán

Quantum Computing and Information Research Group, HUN-REN Wigner RCP



November 27, 2025 AIME 2025, Budapest









Research directions of the group

- Characterization of errors in quantum computers, benchmarking performance
- Quantum machine learning problems (including classical training, but quantum sampling)
- Photonic quantum computing (including ML, and QML)
- Tensor network simulations
- Distillation schemes for quantum networks or distributed quantum computing

Involvement in projects









Benchmarking quantum computers

- What is a **Quantum Computer Benchmark**?
 - Quantum computer benchmarks are structured methods specifically designed to evaluate quantum system performance.
 - A benchmark defines a set of computational tasks for the quantum processor to execute and prescribes how to analyse the resulting data to compute a benchmark score.

Types of benchmarks:

- Component-level benchmarks
 Measure the quality of individual components: gate fidelities, T1/T2 times, measurement errors.
- Device-level benchmarks
 Measure the entire processor's ability to run circuits: volumetric benchmarks, layer fidelities.
- Algorithmic benchmarks
 Compare the performance of entire families of protocols executing algorithms.

Quantum Benchmarks for different QC eras

Era of Quantum Computing	NISQ	late-NISQ	Early Fault-Tolerant	Full Fault-Tolerant
Typical gate errors	10 ⁻¹ - 10 ⁻³ (physical 2Q-gate)	10 ⁻³ -10 ⁻⁵ (physical 2Q-gate)	10 ⁻⁵ - 10 ⁻⁸ (logical gate)	10 ⁻⁸ - (logical gate)
Feasible number of gate operations	10 - 1 000	1 000 - 100 000	1 000 - 1M	>1M
Component-level benchmarks	Very relevant	Very relevant	Very Relevant	Relevant
Device-level benchmarks	Relevant	Very Relevant	Relevant	Relevant
Algorithmic benchmarks	Less relevant	Less relevant	Relevant	Very Relevant

Benchmarking quantum computers (device level)

Near-term benchmarking approach: Instead of full-fledged algorithmic benchmarks, it is more natural to consider benchmarks testing algorithmic primitives, i.e., protocols that appear as subtasks within algorithms.

- Rather than complete algorithmic benchmarks, we focus on benchmarks evaluating algorithmic primitives—smaller-scale protocols and subroutines crucial to quantum algorithms.
- This approach offers clear insight into how close we are to achieving practical quantum utility without requiring the execution of a full-scale quantum algorithm.
- It enables tracking and guiding the evolution of quantum computational platforms toward achieving genuine utility.

Criteria for device-level (volumetric) benchmarks

Benchmarks should satisfy (ideally):

- Well-defined: clear set of rules, reproducible
- Platform independent, hardware-agnostic
 - no native gates, no fixed connectivity
 - preferably the unitary is given as a suitably defined global operation with no particular reference to compilation
- Scalable (wrt the circuit sizes relevant for the era)

Additional considerations:

- Application oriented (relevant for the era)
- Fit to the present devices (e.g., number of shots, software tools)
- Well known, often used by companies, academics
- Randomness involved
- Interesting science involved

Quantum Volume

- Quantum Volume uses random square circuits with all-to-all connectivity
 - d sequential layers acting on d qubits
 - a random permutation of these qubits
 - Haar-random SU(4) operations performed on neighboring pairs

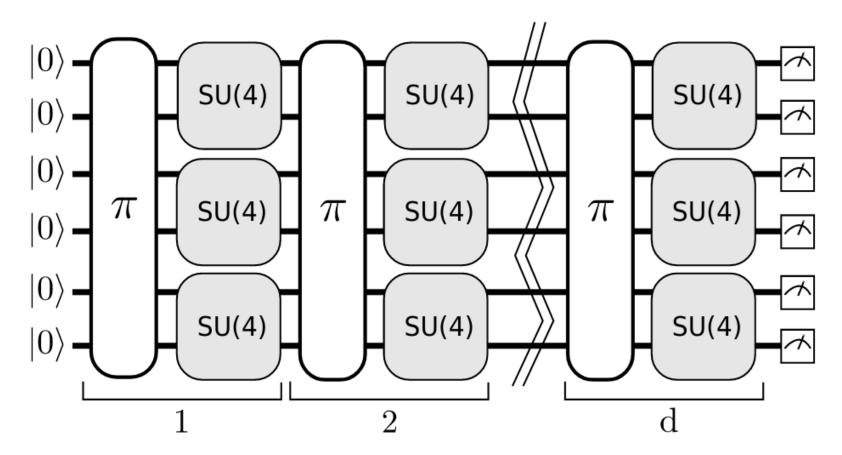
Definition of Quantum Volume

$$\log_2 V_Q = \operatorname*{argmax}_m \min(m, d(m))$$

- m: number of qubits
- d(m): the number of qubits in the largest square circuits that can reliably sample heavy outputs with p > 2/3

Advantages

- It was well tailored for the NISQ era: fit to the present devices, software tools
- Well known, often used by companies, academics
- Interesting theory/physics involved



A. W. Cross, L. S. Bishop, S. Sheldon, P. D. Nation, J. Gambetta, Validating quantum computers using randomized model circuits, Phys. Rev. A 100, 032328 (2019).

Disadvantages

- Not scalable for the late-NISQ (and the Fault Tolerant) era
- The value is not natural (an artificial exponential involved)
- Partially platform biased
- Not application oriented

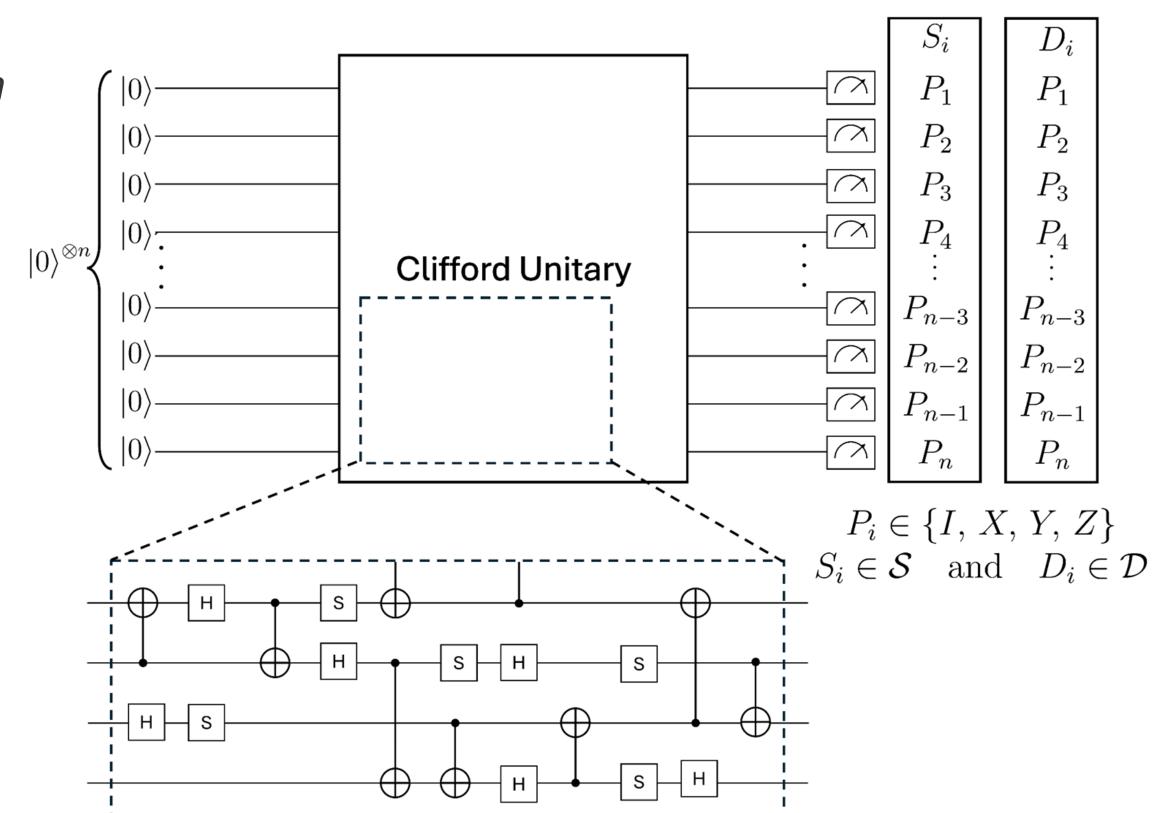
Our volumetric framework

- Computational Task:
 - Algorithmic primitives
 - Quantum operations that remain classically simulable.
 - Trade-off between scalability and quantum utility demonstration.
 - Abstract definition for flexible compilation and optimization, ensuring platform independence.
- Measurement Procedure:
 - Measure specifically chosen expectation values to verify correct quantum state preparation in noise-free conditions.
 - Success criteria
- Benchmark Score:
 - Largest number of qubits on which the quantum processor consistently meets success criteria.

Initialization and Circuit Preparation:

The following steps are carried out by the <u>classical</u> part of the quantum computational stack.

- 1. Randomly select M=10 unitaries from the n-qubit Clifford group.
- 2. For each selected Clifford operations, determine *n* generators of the stabilizer group, additionally, randomly select n distinct *n*-qubit Pauli operators that lie outside the stabilizer group.
- 3. For n > 10, we restrict the set of generators by randomly selecting $10 + \lfloor n/10 \rfloor$ stabilizer operators and an equal number of nonstabilizer operators from the full set.
- 4. Construct and compile the circuit implementation of each n-qubit Clifford, optimised for the benchmarked platform.

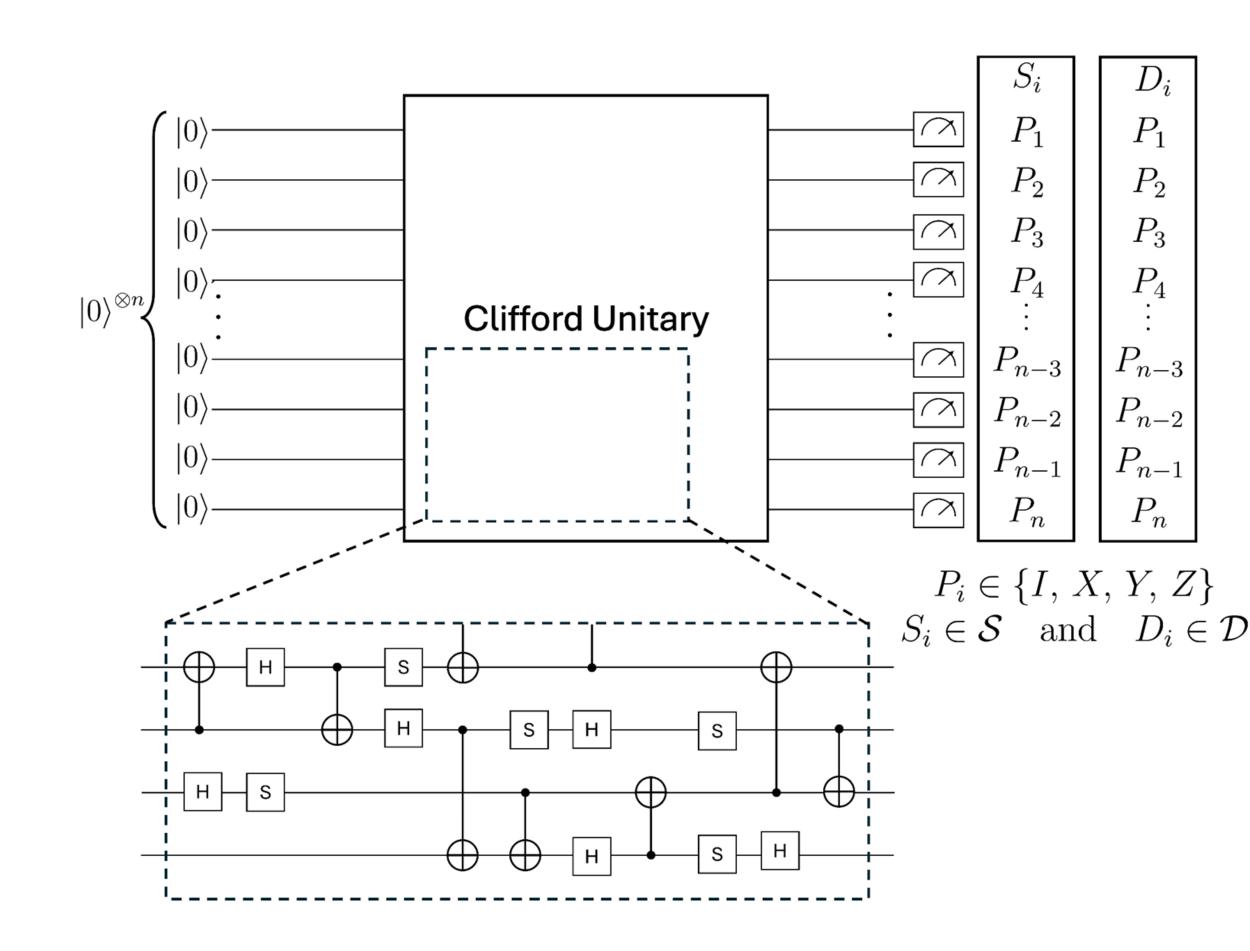


Circuit Execution:

This step must be carried out on a quantum device.

For each Clifford unitary operation

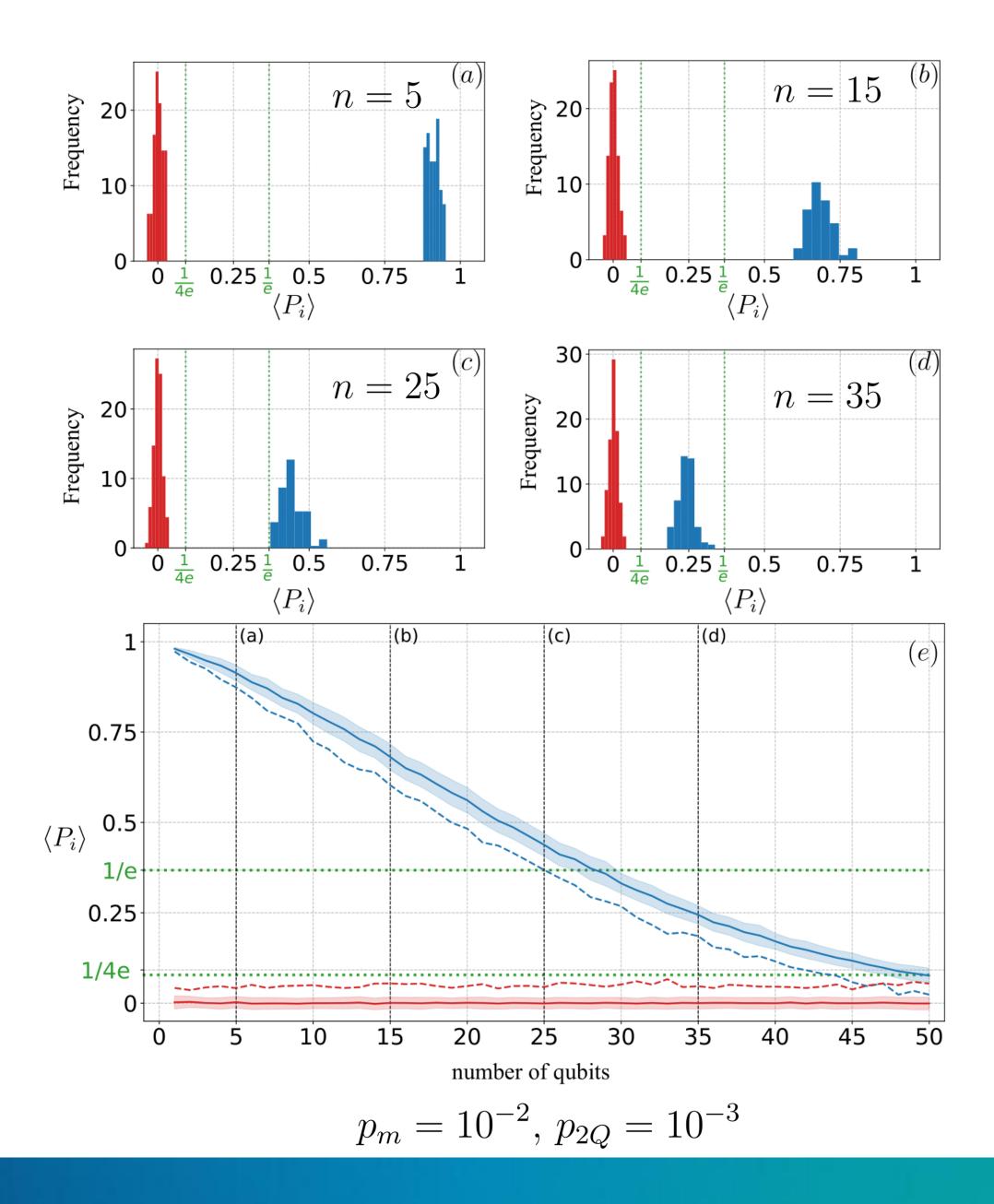
- 1. Initialize all qubits in the $|0\rangle$ state.
- 2. Execute the circuit implementation 2¹² times for each selected **stabilizer** and **non-stabilizer operators**.
- 3. Measure the **expectation value** of the selected **stabilizer** and **non-stabilizer operators.**



Performance Evaluation: We consider the *n*th step of the benchmark successful if the following conditions are fulfilled simultaneously:

$$\min_{m,i} \langle S_i^m \rangle \geq \frac{1}{e} \quad \text{for every chosen Clifford unitary and corresponding stabilizer generator}$$

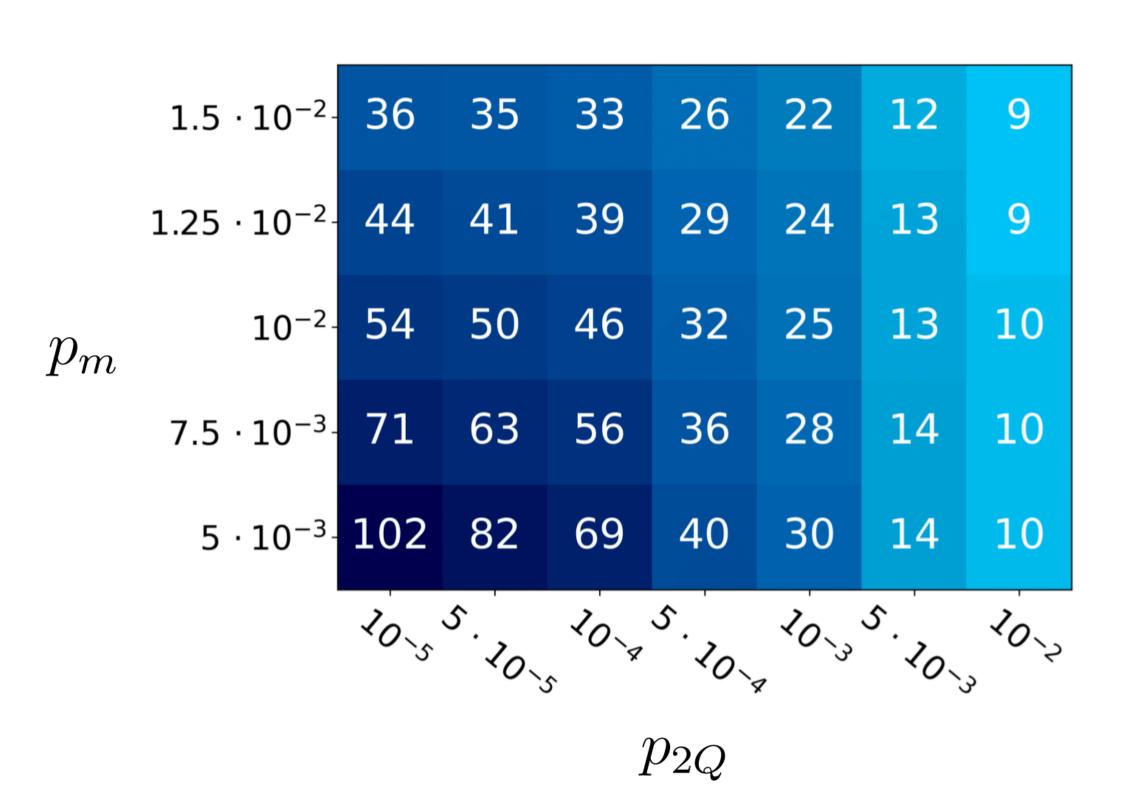
$$\max_{m,i} |\langle D_i^m \rangle| \leq \frac{1}{4e} \quad \text{for every chosen Clifford unitary and corresponding non-stabilizer operator}$$



Performance Evaluation: We consider the *n*th step of the benchmark successful if the following conditions are fulfilled simultaneously:

$$\min_{m,i} \langle S_i^m \rangle \geq \frac{1}{e} \quad \text{for every chosen Clifford unitary and corresponding stabilizer generator}$$

$$\max_{m,i} |\langle D_i^m \rangle| \leq \frac{1}{4e} \quad \text{for every chosen Clifford unitary and corresponding non-stabilizer operator}$$



Advantages:

- Platform-independence: The unitaries are selected from a group of global operations, we avoid prescribing specific compilation implementations.
- Application-oriented: Clifford circuits form the backbone of many fault-tolerant protocols, Shadow tomography, energy estimations,
- Scalability: Classically simulable operations, validates the benchmark results using stabilizer fidelity witnesses, ensuring the benchmark remains scalable.

Disadvantage:

Clifford operations alone cannot achieve quantum advantage, but ...

Free fermionic systems are defined by their quadratic Hamiltonians of the form

$$H=rac{i}{4}\sum_{j,k}A_{j,k}m_jm_k$$
 where A is a real antisymmetric matrix, and m_j are the Majorana operators.

The corresponding unitary operation is given by

$$F(t) = \exp\left(\frac{1}{4} \sum_{i,j=1}^{2d} [\log(O(t))]_{ij} m_i m_j\right)$$
 where $O(t) = e^{tA} \in SO(2L)$

Considering the time evolution governed by the free-fermion operator, we can verify the orthogonality relation for the matrix O by measuring the expected values of the corresponding Majorana mode operators, given by

$$\langle m_j \rangle_{F\rho_i F^{\dagger}} = \text{Tr}\left(F\rho_i F^{\dagger} m_j\right) = \sum_k O_{jk} \langle m_k \rangle_{\rho_i} = \sum_k O_{jk} \delta_{ik} = O_{ji}$$

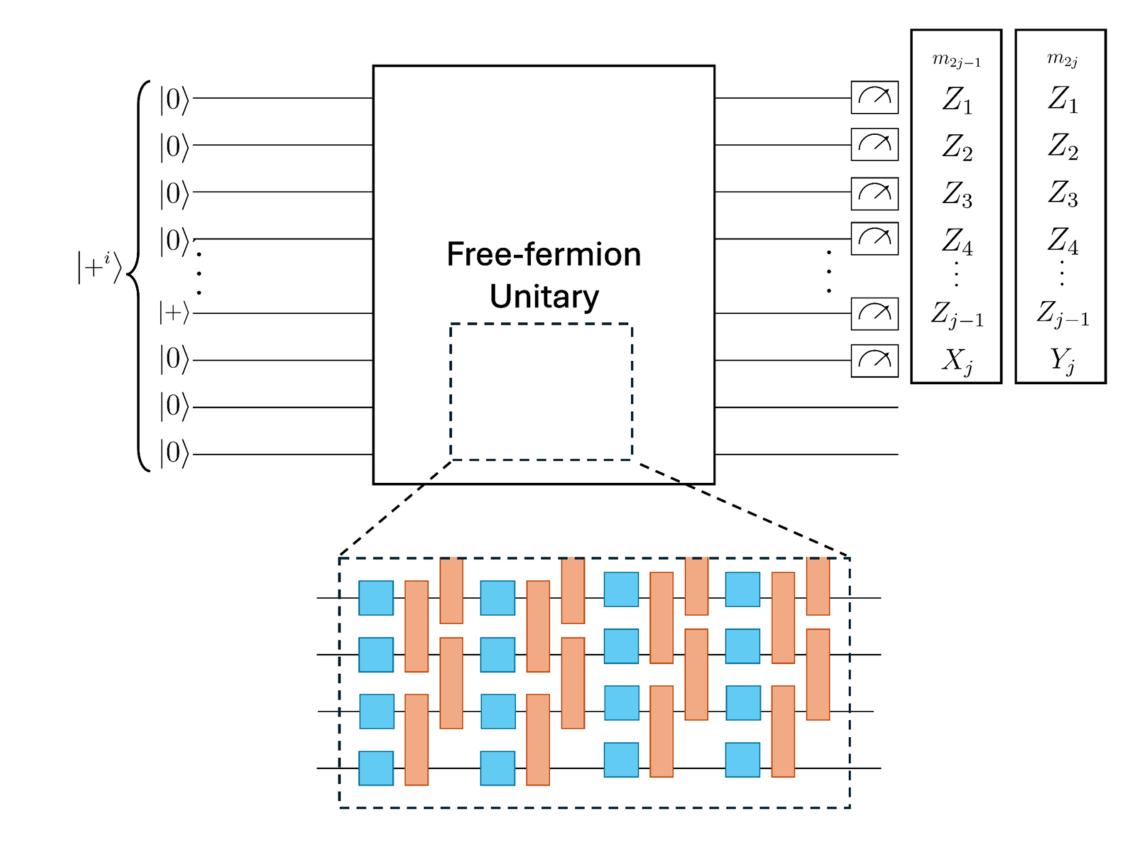
This allows us to test the orthogonality relation by measuring the expected values of the Majorana operators:

$$\sum_{k} O_{kj} \langle m_k \rangle_{F\rho_i F^{\dagger}} = \sum_{k} O_{ki} O_{kj} = \delta_{ij} .$$

Initialization and Circuit Preparation:

The following steps are carried out by the <u>classical</u> part of the quantum computational stack.

- 1. Randomly select M=10 matrices from the SO(2n) group corresponding to n-qubit free-fermion unitaries (matchgates).
- 2. If n > 10, we select the $20 + \lfloor 2n / 10 \rfloor$ largest expectation values determined by the chosen **O matrices** of the SO(2*n*) group, these determine the **Majorana operators** to be measured.
- 3. Construct and compile a circuit implementation of each freefermion unitary.



Circuit Execution:

The following steps must be carried out on a quantum device.

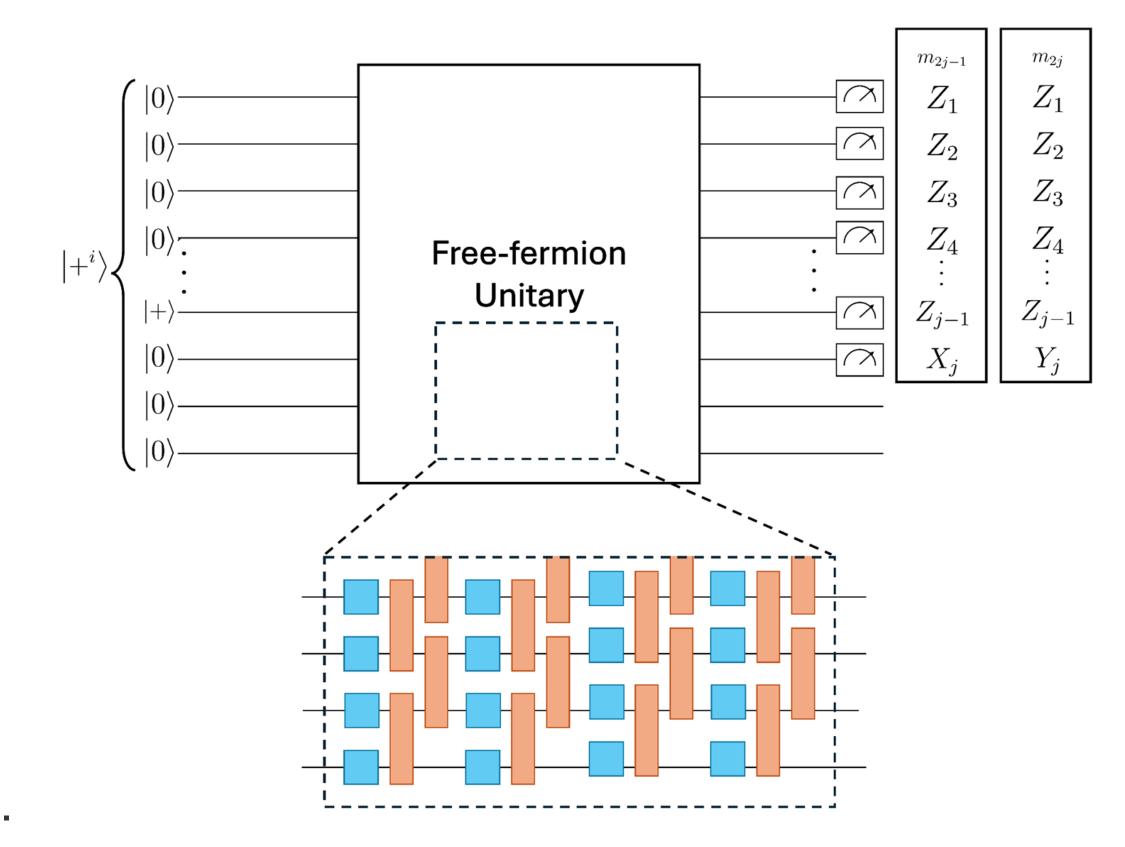
For each Free-fermion unitary

1. Initialize the system in a randomly chosen uniform superposition of fermionic states:

$$\rho_i = \left| +^i \right\rangle \left\langle +^i \right|$$
 where $\left| +^i \right\rangle = \left| 0 \right\rangle^{\otimes i-1} \otimes \left| + \right\rangle \otimes \left| 0 \right\rangle^{\otimes n-i}$

- 2. Execute the circuit implementation.
- 3. Measure the expectation value of each chosen Majorana operators.

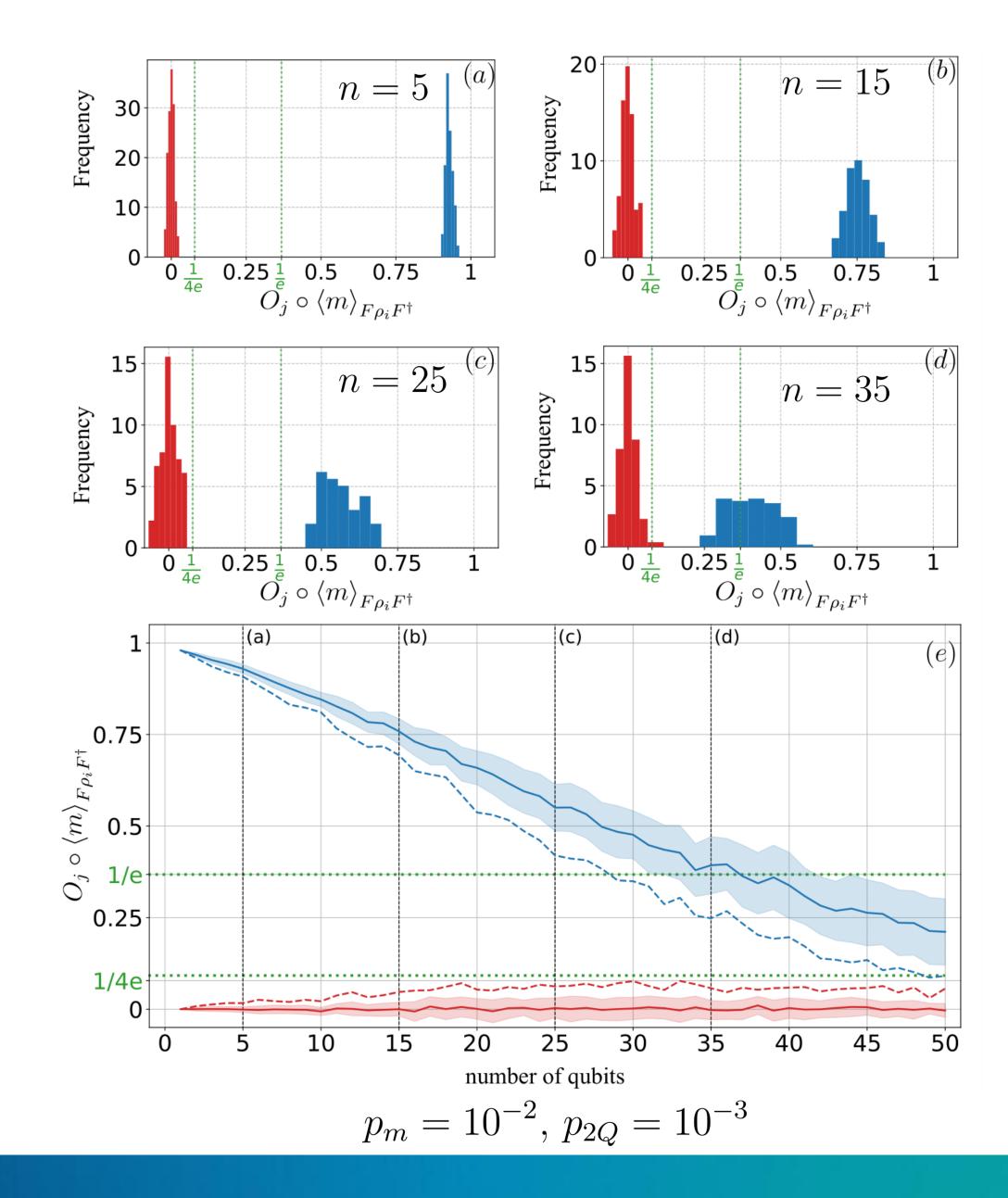
$$m_{2p-1} = Z_1 Z_2 \cdots Z_{p-1} X_p, \quad m_{2p} = Z_1 Z_2 \cdots Z_{p-1} Y_p$$



Performance Evaluation: We consider the *n*th step of the benchmark successful if the following conditions are fulfilled simultaneously:

$$\min_{q} \left(\sum_{k} O_{ki}^{(q)} \left\langle m_{k}^{(q)} \right\rangle \right) \ge \frac{1}{e}$$

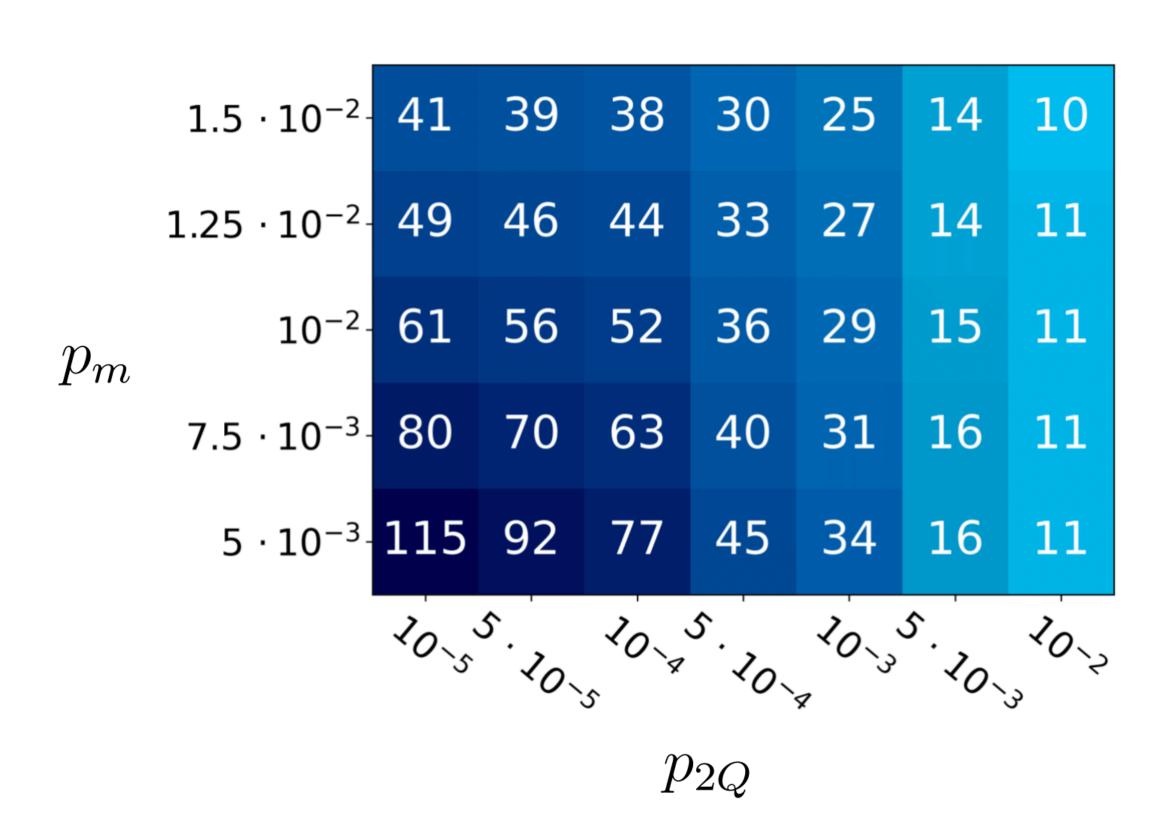
$$\max_{q} \left(\sum_{k} O_{kj}^{(q)} \left\langle m_{k}^{(q)} \right\rangle \right) \le \frac{1}{4e} \quad \text{for j } \neq i$$



Performance Evaluation: We consider the *n*th step of the benchmark successful if the following conditions are fulfilled simultaneously:

$$\min_{q} \left(O^{(q)} \left\langle m_k^{(q)} \right\rangle \right) \ge \frac{1}{e} \qquad \text{for } \mathbf{k} = \mathbf{i}$$

$$\max_{q} \left(O^{(q)} \left\langle m_k^{(q)} \right\rangle \right) \le \frac{1}{4e} \quad \text{for k } \neq \mathbf{i}$$



Advantages:

- Platform-independence: The unitaries are selected from a group of global operations, we avoid prescribing specific compilation implementations.
- **Application-oriented**: energy estimations, fermion-to-qubit mapping changes, simulation of free-fermion Hamiltonians, quantum chemistry, quantum many-body models
- Scalability: Classically simulable operations, validates the benchmark results using fidelity witnesses, ensuring the benchmark remains scalable.
- Together with Clifford unitaries they form a universal gate-set, thus when applied jointly they are hard to simulate classically. Example: Fermion Sampling, which can be validated by measuring the 4-point correlators while sampling from the output state is classically hard.

M. Oszmaniec, N. Dangniam, M. E. S. Morales, Z. Zimborás, Fermion sampling: a robust quantum computational advantage scheme using fermionic linear optics and magic input states, PRX Quantum, 2022.

From the European Quantum Flagship's KPI Booklet

Key Performance Indicators for Quantum Computing

- 1. Circuit size: The largest random N-qubit Clifford operation that can be reliably executed
- 2. Multipartite Entanglement: The largest GHZ state successfully prepared and verified by a quantum computer
- 3. Cryptanalysis: Largest reliable instance of Shor's algorithm.
- **4. Error correction:** the ratio of the best achievable Bell state fidelity on physical qubits to the best achievable Bell state fidelity on logical qubits on the same hardware





