## Entanglement and Codes

C. Eltschka, O. Gühne, M. Grassl, F. Huber, J. Siewert

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## Overview

Part I: Absolutely maximally entangled states
Part II: Entanglement and codes
Part III: Highly entangled subspaces / QMDS codes

## Part I: Absolutely maximally entangled states



## Motivation

## How entangled can two couples get?

A. Higuchi, A. Sudbery *<br>Dept. of Mathematics, University of York, Heslington, York, YO10 5DD, UK<br>Received 9 June 2000; accepted 12 July 2000<br>Communicated by P.R. Holland

A. Higuchi and A. Sudbery, Phys. Lett. A 273, 213 (2000)

## Absolutely maximally entangled states

Definition (AME states)
A pure state $\left|\phi_{n, D}\right\rangle$ is called absolutely maximally entangled (AME), if it shows maximal entanglement over all bipartitions.
( $\equiv$ all its reductions to $\left\lfloor\frac{n}{2}\right\rfloor$ parties are maximally mixed)

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\text { ( } \equiv \text { all its reductions to }\left\lfloor\frac{n}{2}\right\rfloor \text { parties are maximally mixed) }
$$

## Example

For all prime dimensions (graph states):

$\longrightarrow$ arbitrary dimensions: prime-decomposition $D=p_{1} p_{2} \ldots p_{r}$ :

$$
\left|\phi_{n, D}\right\rangle=\left|\phi_{n, p_{1}}\right\rangle \otimes\left|\phi_{n, p_{2}}\right\rangle \otimes \ldots \otimes\left|\phi_{n, p_{r}}\right\rangle
$$

## Bounds on AME state existence

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Necessary condition for existence:

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n \leq \begin{cases}2\left(D^{2}-1\right) & n \text { even } \\ 2 D(D+1)-1 & n \text { odd }\end{cases}
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- Open problem: provide tight bounds for the existence of AME states (and quantum codes).


## $\nexists$ Four-qubit AME

"Bloch-style" proof by contradiction:
Assume a 4-qubit AME state $\varrho_{A B C D}=|\phi\rangle\langle\phi|$ exists.

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a) From Schmidt decomposition, "projector relation" holds:

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b) Decompose

$$
\varrho_{A B C}=\frac{1}{2^{3}}(\mathbb{1}+\underbrace{\sum_{\alpha, \beta, \gamma \in\{x, y, z\}} c_{\alpha \beta \gamma} \sigma_{\alpha} \otimes \sigma_{\beta} \otimes \sigma_{\gamma}}_{P_{3}(\neq 0)})
$$

Note that there are no terms of e.g. the form

$$
\sum_{\alpha, \beta \in\{x, y, z\}} c_{\alpha \beta} \sigma_{\alpha} \otimes \sigma_{\beta} \otimes \sigma_{0} \quad!
$$

## $\nexists$ Four-qubit AME, continued (I)

c) From projector relation $\varrho_{A B C}^{2}=\frac{1}{2} \varrho_{A B C}$,

$$
\left(P_{3}\right)^{2}=\frac{1}{2}\left\{P_{3}, P_{3}\right\} \stackrel{!}{=} 3 \mathbb{1}+2 P_{3}
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d) Above, in $\left\{P_{3}, P_{3}\right\}$ two different Paulis either

$$
\begin{aligned}
\sigma_{j} \sigma_{k} & =i \epsilon_{j k l} \sigma_{l}, \quad j \neq k \\
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e) To contribute to $P_{3}$ on the RHS, three pairs of Paulis need to produce three other Paulis. Factor of $i^{3}$ appears, and term vanishes in anticommutator. Thus $P_{3}=0$. Contradiction!
$\Longrightarrow \nexists$ four-qubit AME state.

## $\nexists$ Seven-qubit AME

Proof by contradiction:
Assume a 7-qubit AME state $\varrho=|\phi\rangle\langle\phi|$ exists.
(a) We use the Bloch decomposition and sort the correlations:

$$
\varrho \sim \sum_{\alpha_{1} \ldots \alpha_{n}} r_{\alpha_{1}, \ldots, \alpha_{n}} \sigma_{\alpha_{1}} \otimes \cdots \otimes \sigma_{\alpha_{N}} \sim\left(\mathbb{1}^{\otimes n}+\sum_{j=4}^{7} P_{j}\right)
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(b) General (qubit) parity rule for $\left\{P_{j}, P_{k}\right\}$ :

$$
\begin{array}{rll}
\text { \{even, even }\} & \longrightarrow & \text { even } \\
\{\text { odd, odd }\} & \longrightarrow & \text { even } \\
\{\text { odd, even }\} & \longrightarrow & \text { odd } .
\end{array}
$$

## $\nexists$ Seven-qubit AME, continued (I)

(c) The four- and five-qubit reductions fulfill "projector relations"

$$
\varrho_{(4)}^{2}=\frac{1}{8} \varrho_{(4)} \quad \varrho_{(5)}^{2}=\frac{1}{4} \varrho_{(5)} .
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$$

(d) Expand $\varrho_{(4)}$ and $\varrho_{(5)}$ in the Bloch basis

$$
\varrho_{(4)}=\frac{1}{2^{4}}\left(\mathbb{1}+P_{4}\right), \quad \varrho_{(5)}=\frac{1}{2^{5}}\left(\mathbb{1}+\sum_{j=1}^{5} P_{4}^{[j]} \otimes \mathbb{1}^{(j)}+P_{5}\right) .
$$

## $\nexists$ Seven-qubit AME, continued (II)

(e) Resulting eigenvalue equations:

$$
P_{4}^{[j]} \otimes \mathbb{1}^{\otimes 3}|\phi\rangle=1|\phi\rangle, \quad P_{5} \otimes \mathbb{1}^{\otimes 2}|\phi\rangle=2|\phi\rangle .
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(f) Expanding $\varrho_{(5)}^{2}=\frac{1}{4} \varrho_{(5)}$ gives two equations (parity rule).

$$
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$$

(g) Multiplying with $|\phi\rangle$ from the right:

$$
(5 \cdot 1 \cdot 2+2 \cdot 5 \cdot 1)|\phi\rangle \neq 6 \cdot 2|\phi\rangle .
$$

$\Longrightarrow \nexists$ seven-qubit AME. (similar contradiction found for all $n \neq 2,3,5,6$.)

## A best approximation...

## Result

A seven qubit AME does not exist. At most 32 out of 35 three-body RDMs can be maximally mixed.


FH, O. Gühne, J. Siewert, Phys. Rev. Lett. 118, 200502 (2017)

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- Consider correlation constraints from generalized state inversion / shadow inequality (talk by Jens)

$$
\operatorname{Tr}\left(\mathcal{I}_{T}[\varrho] \varrho\right)=\sum_{S \subseteq\{1 \ldots n\}}(-1)^{|S \cap T|} \operatorname{tr}\left[\varrho_{S}^{2}\right] \geq 0
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Example ( $\nexists$ four-qubit AME)

$$
\begin{aligned}
\operatorname{Tr}\left(\mathcal{I}_{1234}[\varrho] \varrho\right) & =1-\sum_{i} \operatorname{tr}\left(\rho_{i}^{2}\right)+\sum_{i<j} \operatorname{tr}\left(\varrho_{i j}^{2}\right)-\sum_{i<j<k} \operatorname{tr}\left(\rho_{i j k}^{2}\right)+\operatorname{tr}\left(\varrho^{2}\right) \\
& =1-4 \frac{1}{2}+6 \frac{1}{4}-4 \frac{1}{2}+1=-\frac{1}{2} \nsucceq 0
\end{aligned}
$$

## Bounds on higher-dimensional AME states

## Further bounds

A further 27 higher-dimensional AME states $\nexists$ (light blue).


$\exists$ : state exists
dark blue: excluded by Scott's bound
light blue: excluded by the shadow inequality

## Mixed-dimensional AME states

Consider maximally entangled systems of mixed dimensions (e.g. qubit-qutrit), with maximal entanglement across every bipartition:
$2 \times 2 \times 2 \times 2$ : $\nexists$ four-qubit AME (proof at the beginning)
$2 \times 2 \times 2 \times 3$ : $\nexists$ shadow inequality
$2 \times 2 \times 3 \times 3$ : $\nexists$ shadow inequality
$2 \times 3 \times 3 \times 3$ : $\exists$ see new state below
$3 \times 3 \times 3 \times 3$ : $\exists$ four-qutrit AME (c.f. Karol's talk)

$$
\begin{aligned}
\left|\phi_{2333}\right\rangle= & -\alpha|0011\rangle-\beta|0012\rangle+\beta|0021\rangle+\alpha|0022\rangle \\
& -\beta|0101\rangle+\alpha|0102\rangle+\beta|0110\rangle+\alpha|0120\rangle \\
& -\alpha|0201\rangle+\beta|0202\rangle-\alpha|0210\rangle-\beta|0220\rangle \\
& -\beta|1011\rangle+\alpha|1012\rangle-\alpha|1021\rangle+\beta|1022\rangle \\
& +\alpha|1101\rangle+\beta|1102\rangle-\alpha|1110\rangle+\beta|1120\rangle \\
& -\beta|1201\rangle-\alpha|1202\rangle-\beta|1210\rangle+\alpha|1220\rangle \\
& 12\left(\alpha^{2}+\beta^{2}\right)=1,54 \alpha \beta=1 .
\end{aligned}
$$

FH, C. Eltschka, J. Siewert, O. Gühne, J. Phys. A: Math. Theor. 51, 175301 (2018)

## Part II: Entanglement and Quantum Codes



## Quantum codes

A quantum code is a subspace of a multipartite system: Denote by $\mathcal{Q}$ a subspace of $\left(\mathbb{C}^{d}\right)^{\otimes n}$ spanned by an ONB $\left\{\left|v_{i}\right\rangle\right\}$. Let $\Pi=\sum_{i}^{K}\left|v_{i}\right\rangle\left\langle v_{i}\right|$ be the projector onto it, with $\operatorname{rank}(\Pi)=K$.

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## Theorem (Knill-Laflamme error-conditions)

The subspace $\mathcal{Q}$ is a QECC of distance at least $d$, if and only if for all operators with $|\operatorname{supp}(E)|<d$,

$$
\left\langle v_{i}\right| E\left|v_{j}\right\rangle=\delta_{i j} C_{E} \quad\left(=\delta_{i j} \operatorname{tr}[E] \text { "pure" }\right)
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$$

- If distance is $d$, then errors on $\left\lfloor\frac{(d-1)}{2}\right\rfloor$ particles can be corrected.
- $\mathcal{Q}$ is denoted as a $((n, K, d))_{D}$ code.

$$
\text { E. Knill, R. Laflamme, and L. Viola, Phys. Rev. Let. 84, } 2525 \text { (2000). }
$$

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a) For all $|\phi\rangle \in \mathcal{Q}$, and all subsets $|S|<d$

$$
\operatorname{tr}_{s c}(|\phi\rangle\langle\phi|)=\varrho_{s} \quad\left(=\mathbb{1} / D^{|S|}\right)
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\operatorname{tr}_{S^{c}}(|\phi\rangle\langle\phi|)=\varrho_{S} \quad\left(=\mathbb{1} / D^{|S|}\right)
$$

$\longrightarrow$ "every vector looks locally the same"
b) Let $\varrho=\Pi / K$. For all subsets $|S|<d$,

$$
K \operatorname{tr}\left[\varrho_{S^{c}}^{2}\right]=\operatorname{tr}\left[\varrho_{S}^{2}\right] \quad\left(=1 / D^{|S|}\right)
$$

$\longrightarrow$ "constraints on purities of complementary reductions"
E. Rains, IEEE Trans. Inf. Theory 44, 4 (1998)

## Part III: QMDS codes \& highly entangled subspaces

## Highly entangled subspaces

## Definition

A pure state $|\phi\rangle$, whose reductions onto $r$ parties are all maximally mixed, is termed $r$-uniform. A $r$-uniform subspace (rUS) is a subspace of $\left(\mathbb{C}^{D}\right)^{\otimes n}$, in which every vector is at least $r$-uniform.

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## Observation (pure QECC $\equiv \mathrm{r}$-uniform subspace)

The following objects are equivalent:
a) a pure $((n, K, d))_{D}$ quantum error correcting code.
b) a $(d-1)$-uniform subspace in $\left(\mathbb{C}^{D}\right)^{\otimes n}$ of dimension $K$.

## Bounds on codes

Theorem (Quantum Singleton bound)
Let $\mathcal{Q}$ be a $((n, K, d))_{D}$ quantum error correcting code. Then

$$
n+2 \geq \log _{D} K+2 d
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- If equality above, the code is called quantum maximum distance separable (QMDS)
- Fact: QMDS codes are pure (have maximally mixed marginals).
$\longrightarrow$ QMDS codes are the largest possible $r$-uniform subspaces.


## New codes from old

New codes can be constructed from old ones:
Theorem
Let $((n, K, d))_{D}$ be a pure QECC with $n, d \geq 2$. Then there exists a pure code $((n-1, D K, d-1))_{D}$.
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$\longrightarrow$ corresponds to taking a partial trace over one particle.

## QMDS families

...apply to QMDS codes:
Example

$$
\begin{array}{l|l}
\left(\left(6,2^{0}, 4\right)\right)_{2} \exists & \left(\left(12,3^{0}, 7\right)\right)_{3} \nexists \\
\left(\left(5,2^{1}, 3\right)\right)_{2} \exists & \left(\left(11,3^{1}, 6\right)\right)_{3} \nexists \\
\left(\left(4,2^{2}, 2\right)\right)_{2} \exists & \left(\left(10,3^{2}, 5\right)\right)_{3} \nexists \\
\left(\left(3,2^{3}, 1\right)\right)_{2} \exists & \left(\left(9,3^{3}, 4\right)\right)_{3} \\
& \nexists \\
& \left(\left(8,3^{4}, 3\right)\right)_{3} \\
& \exists \\
& \left(\left(7,3^{5}, 2\right)\right)_{3} \\
& \exists \\
& \left(\left(6,3^{6}, 1\right)\right)_{3} \\
\exists
\end{array}
$$

- Family of codes / highly entangled subspaces determined by $n+k$.
- For a given family, if the parent-AME does not exist, what is the uppermost member?


## Bound on the existence of QMDS codes

## Maximal length of QMDS codes

$\mathrm{A}((n, K, d))_{D}$ QMDS code of distance $d \geq 3$
$\left[\equiv(d-1)\right.$-uniform subspace in $\left(\mathbb{C}^{D}\right)^{\otimes n}$ of dimension $\left.K\right]$ must satisfy

$$
\begin{aligned}
n & \leq D^{2}+d-2, \quad \text { or equivalently } \\
n+k & \leq 2\left(D^{2}-1\right) .
\end{aligned}
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FH and M. Grassl, in preparation.

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- Extends Scott's AME bound and stabilizer QMDS bounds to all QMDS codes.
- Further bounds from the shadow inequality / generalized inversion $\operatorname{tr}\left(\mathcal{I}_{T}[\varrho\rfloor \varrho\right) \geq 0$.


## Examples

Example
All QMDS－families of local dimension $D=3$ ：

| n＋k | bound | achieved |  |
| :---: | :---: | :---: | :---: |
| 4 | ［4， $0,3 \rrbracket_{3}$ | 【4， $0,3 \rrbracket_{3}$ | （optimal） |
| 6 | ［6， $0,4 \rrbracket_{3}$ | ［6， $0,4 \rrbracket_{3}$ | （optimal） |
| 8 | ［6， $2,3 \rrbracket_{3}$ | ［6， $2,3 \rrbracket_{3}$ | （optimal） |
| 10 | 【10， $0,6 \rrbracket_{3}$ | 【10， $0,6 \rrbracket_{3}$ | （optimal） |
| 12 | ［ $8,4,3 \rrbracket_{3}$ | 【8， $4,3 \rrbracket_{3}$ | （optimal） |
| 14 | 【11， $3,5 \rrbracket_{3}$ | $\llbracket 10,4,4 \rrbracket_{3}$ |  |
| 16 | $\llbracket 11,5,4 \rrbracket_{3}$ | $\llbracket 10,6,3 \rrbracket_{3}$ |  |

## Summary of Results

- Arbitrarily strong quantum correlations are not allowed. Qubit AME states only exist for $n=2,3,5,6$.



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- Bounds for the existence QMDS codes / highly entangled subspaces

FH and M. Grassl, in preparation.

## Thank you for your attention ...

... and thanks to my collaborators!


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Swiss National Science Foundation DFG CEELEX

## IR

Universität Regensburg

